CE 815 – Secure Software Systems

ML-Based Vulnerability Detection Methods (GraphSPD)

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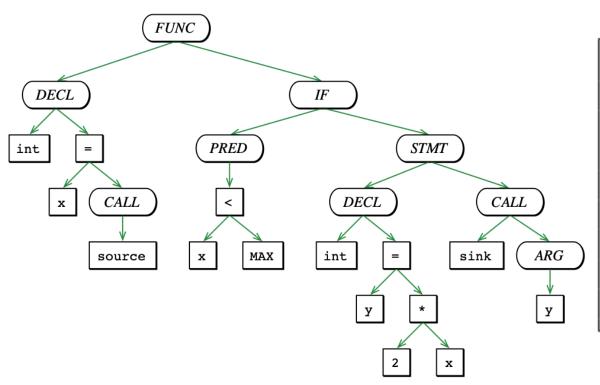
Acknowledgments: Some of the slides are fully or partially obtained from other sources. A reference is noted on the bottom of each slide, when the content is fully obtained from another source. Otherwise a full list of references is provided on the last slide. Thanks to Mohammad Haddadian for the help on the slides.



```
void foo()
{
   int x = source();
   if (x < MAX)
   {
      int y = 2 * x;
      sink(y);
   }
}</pre>
```

Abstract Syntax Tree





```
void foo()
{
    int x = source();
    if (x < MAX)
    {
        int y = 2 * x;
        sink(y);
    }
}</pre>
```

Control Flow Graph

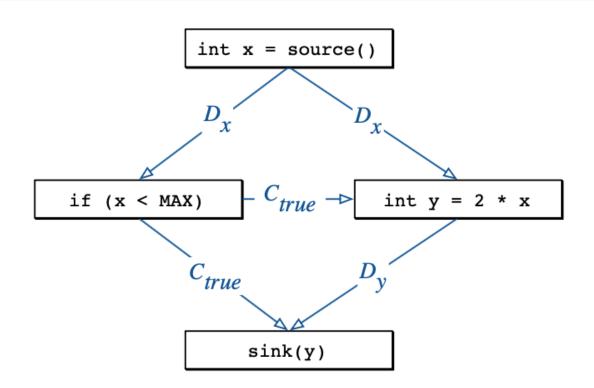


```
ENTRY
int x = source()
  if (x < MAX)
         true
 int y = 2 * x
                      false
    sink(y)
      EXIT
```

```
void foo()
{
   int x = source();
   if (x < MAX)
   {
      int y = 2 * x;
      sink(y);
   }
}</pre>
```

Program Dependence Graph

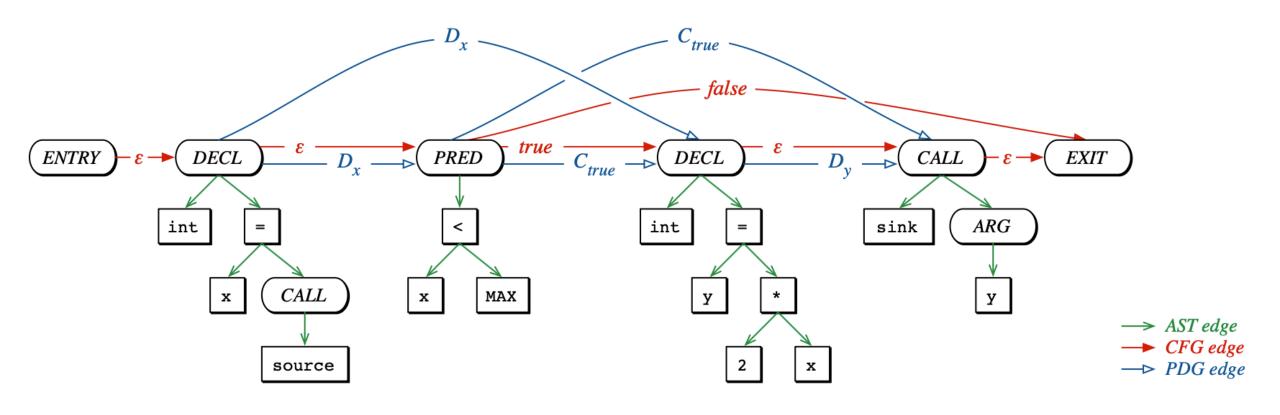




```
void foo()
{
    int x = source();
    if (x < MAX)
    {
        int y = 2 * x;
        sink(y);
    }
}</pre>
```

Code Propert Graph







GraphSPD: Graph-Based Security Patch Detection with Enriched Code Semantics, S. Wang, X. Wang, K. Sun, S. Jajodia, H. Wang, Q. Li, IEEE S&P 2023.

Introduction



- As announced by the 2021 report, 98% of codebases contain open source components
- Meanwhile, 84% of code-bases have at least one open-source vulnerability
- 60% of them contain high-risk vulnerabilities

 By exploiting the OSS vulnerabilities reported in the vulnerability databases (e.g., NVD), attackers can perform "N-day" attacks against unpatched software systems

Problem



• A large volume of OSS security patches (e.g., GitHub commits fixing vulnerabilities) are silently released.

From 7f9822a48213dd2feca845dbbb6bcb8beb9550de Subject: [PATCH] Add blinding to a DSA signature

This is based on side channel attacks demonstrated by (NCC Group) for ECDSA which are likely to be able to be applied to DSA.

Not report to NVD

From 41bdc78544b8a93a9c6814b8bbbfef966272abbe Subject: [PATCH] x86/tls: Validate TLS entries to protect espfix

Installing a 16-bit RW data segment into the GDT defeats espfix. AFAICT this will not affect glibc, Wine, or dosemu at all.

Not provide explicit description

 Average users need to timely detect and apply security patches before being exploited by armored attackers.

Previous Solutions and Limitations



An OSS Patch

Natural Language Documentation (Commit Message/Changelog)

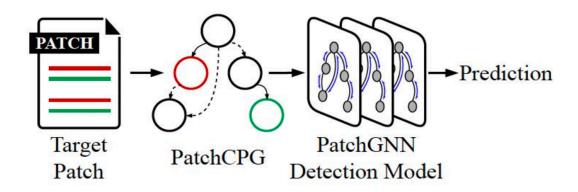
Source Code Changes

- Mining security keywords
 - Requiring well-maintained doc.

- Regarding code as sequential data
 - Losing important semantics.
- Our solution: representing code as graph
 - Retaining rich patch structural info.

A Graph-Based Security Patch Detection System



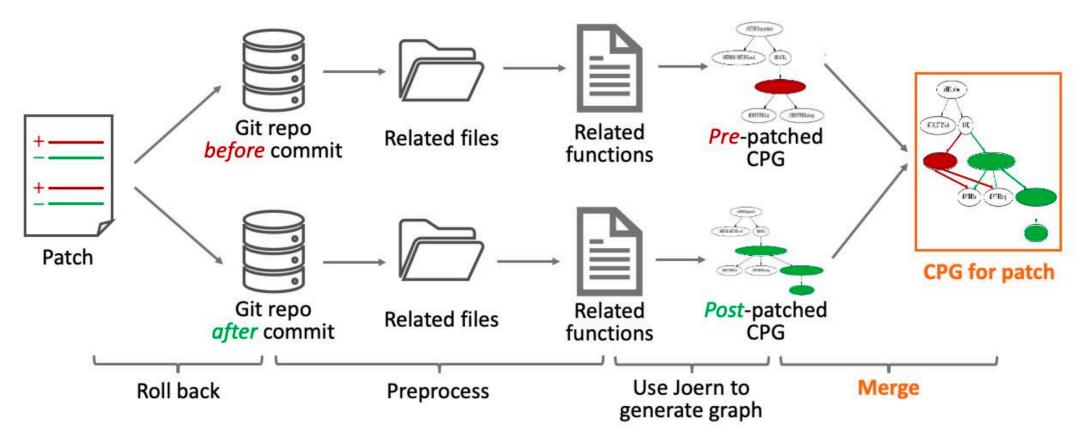


- PatchCPG: a new graph representation of inherent code change structures.
 - Syntax and semantics: AST + control & data dependency graph.
 - Changes and relations with context: pre-patch + post-patch graph.
- PatchGNN: a tailored GNN model to capture diverse patch structural information

PatchCPG: From Patch to Graph

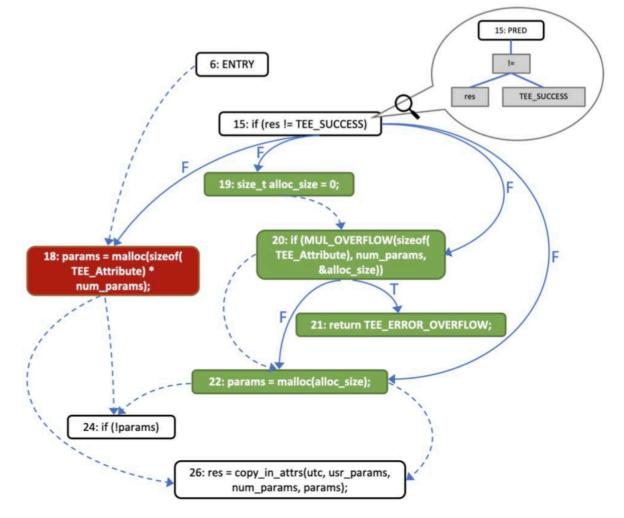


Challenge: how to construct PatchCPG?





 A joint graph encodes rich patch structural information.





• A joint graph encodes rich 15: PRED patch structural information. 6: ENTRY TEE SUCCESS 15: if (res != TEE SUCCESS) 19: size_t alloc_size = 0; 20: if (MUL OVERFLOW(sizeof(TEE_Attribute), num_params, 18: params = malloc(sizeof(What is deleted? TEE Attribute) * &alloc size)) num params); 21: return TEE ERROR OVERFLOW; What is added? 22: params = malloc(alloc_size) 24: if (!params) 26: res = copy in attrs(utc, usr params, What context statements are related? num_params, params);



 A joint graph encodes rich 15: PRED patch structural information. 6: ENTRY TEE_SUCCESS 15: if (res != TEE_SUCCESS) 19: size_t alloc_size = 0; Which statements decide the (un)safe operation? 20: if (MUL_OVERFLOW(sizeof(TEE_Attribute), num_params, 18: params = malloc(sizeof(&alloc_size)) TEE_Attribute) * What is deleted? num_params); 21: return TEE ERROR OVERFLOW: What is added? 22: params = malloc(alloc_size) 24: if (!params) 26: res = copy_in_attrs(utc, usr_params, num params, params); What context statements are related?



 A joint graph encodes rich 15: PRED patch structural information. 6: ENTRY TEE SUCCESS 15: if (res != TEE_SUCCESS) 19: size_t alloc_size = 0; Which statements decide the (un)safe operation? 20: if (MUL_OVERFLOW(sizeof(TEE_Attribute), num_params, 18: params = malloc(sizeof(What is deleted? TEE_Attribute) * &alloc_size)) num_params); 21: return TEE ERROR OVERFLOW; What is added? Where the value comes from? 22: params = malloc(alloc_size) 24: if (!params) 26: res = copy_in_attrs(utc, usr_params, num params, params); What context statements are related?



How each statement looks like? A joint graph encodes rich 6: ENTRY patch structural information. TEE_SUCCESS 15: if (res != TEE_SUCCESS) 19: size_t alloc_size = 0; Which statements decide the (un)safe operation? 20: if (MUL_OVERFLOW(sizeof(TEE Attribute), num params, 18: params = malloc(sizeof(&alloc_size) TEE Attribute) * What is deleted? num params): 21: return TEE_ERROR_OVERFLOW; What is added? 22: params = malloc(alloc_size); Where the value comes from? 24: if (!params) 26: res = copy in attrs(utc, usr params, num_params, params); What context statements are related?

Reducing Noisy Information by Slicing

slicing

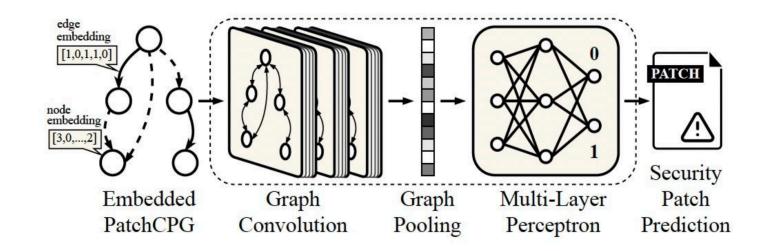


```
6 TEE_Result syscall_asymm_verify(unsigned long state,
       const struct utee_attribute *usr_params, size_t
       num_params, const void *data, size_t data_len,
       const void *sig, size_t sig_len)
       TEE Result res;
       TEE_Attribute *params = NULL;
       struct user_ta_ctx *utc;
11
12
                         Too many statements
13
       res = tee_mmu_check_access_rights(utc,
       TEE_MEMORY_ACCESS_READ
       TEE_MEMORY_ACCESS_ANY_OWNER, (uaddr_t) sig, sig_len)
15
       if (res != TEE SUCCESS)
           return res;
      params = malloc(sizeof(TEE_Attribute) * num_params);
      size_t alloc_size = 0;
      if (MUL_OVERFLOW(sizeof(TEE Attribute), num params,
       &alloc size))
21 +
            return TEE_ERROR_OVERFLOW;
22 +
      params = malloc(alloc_size);
      if (!params)
           return TEE_ERROR_OUT_OF_MEMORY;
26
       res = copy_in_attrs(utc, usr_params, num_params,
       params);
       if (res != TEE_SUCCESS)
28
           goto out;
29
30
32 out:
33
       free (params);
34
       return res;
35
```

```
6 TEE Result syscall asymm verify (unsigned long state,
       const struct utee_attribute *usr_params, size_t
       num_params, const void *data, size_t data_len,
       const void *sig, size t sig len)
      TEE Result res;
      TEE Attribute *params = NULL;
      struct user_ta_ctx *utc;
      . . .
      res = tee_mmu_check_access_rights(utc,
       TEE MEMORY ACCESS READ
       TEE_MEMORY_ACCESS_ANY_OWNER, (uaddr_t) sig, sig_len)
      if (res != TEE_SUCCESS)
          return res;
      params = malloc(sizeof(TEE Attribute) * num params);
      size_t alloc_size = 0;
      if (MUL_OVERFLOW(sizeof(TEE_Attribute), num_params,
       &alloc size))
           return TEE ERROR OVERFLOW;
      params = malloc(alloc_size);
      if (!params)
          return TEE_ERROR_OUT_OF_MEMORY;
      res = copy_in_attrs(utc, usr_params, num_params,
       params);
      if (res != TEE_SUCCESS)
          goto out;
29
      Only retain most relevant contexts
32 out:
33
      free (params);
34
      return res;
35
```

PatchGNN: Detect Security Patches from PatchCPGs





- Challenge 1: how to embed the PatchCPGs?
- Challenge 2: how to learn multiple attributes (CDG/DDG/AST/pre/post)?

PatchCPG Embeddings



Node Embedding

- 20-dimensional vulnerability features.
 - Code Snippet Metadata (2): the number of characters, the version information (i.e., deleted, added, or context).
 - Identifier and Literal Features (7): the number of function calls, variables, numeric numbers, strings, pointers, arrays, and null identifiers.
 - Control Flow Features (3): the boolean features indicating if the node is a conditional, loop, or jump statement.
 - Operator Features (4): the number of arithmetic, relational, logical, and bitwise operators.
 - API Features (4): the boolean features indicating if the code snippet contains the APIs of memory operations, string operations, lock operations, and system operations.





TABLE I: The involved tokens or sub-tokens of the control flow features, the operator features, and the API features.

Features	Matched Tokens or Sub-tokens			
condition	if, switch			
loop	for, while			
jump	return, break, continue, goto, throw, assert			
arithmetic [†]	++,, +, -, *, /, %, =, +=, -=, *=, /=, %=			
relational	==, !=, >=, <=, >, <			
logical	&&, , !, and, or, not			
bitwise*	&, $, <<, >>, \sim$, \wedge , bitand, bitor, oxr			
memory API	alloc, free, mem, copy, new, open, close, delete, create, release, sizeof, remove, clear, dequene, enquene, detach, attach			
string API	str, string			
lock API	lock, mutex, spin			
system API	init, register, disable, enable, put, get, up, down, inc, dec, add, sub, set, map, stop, start, prepare, suspend, resume, connect,			

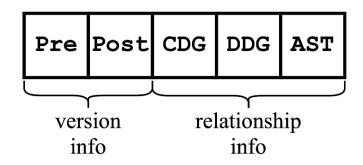
[†] Operator * is determined as dereference operator or arithmetic operator.

^{*} Operator & is determined as address-of operator or bitwise operator.

PatchCPG Embeddings



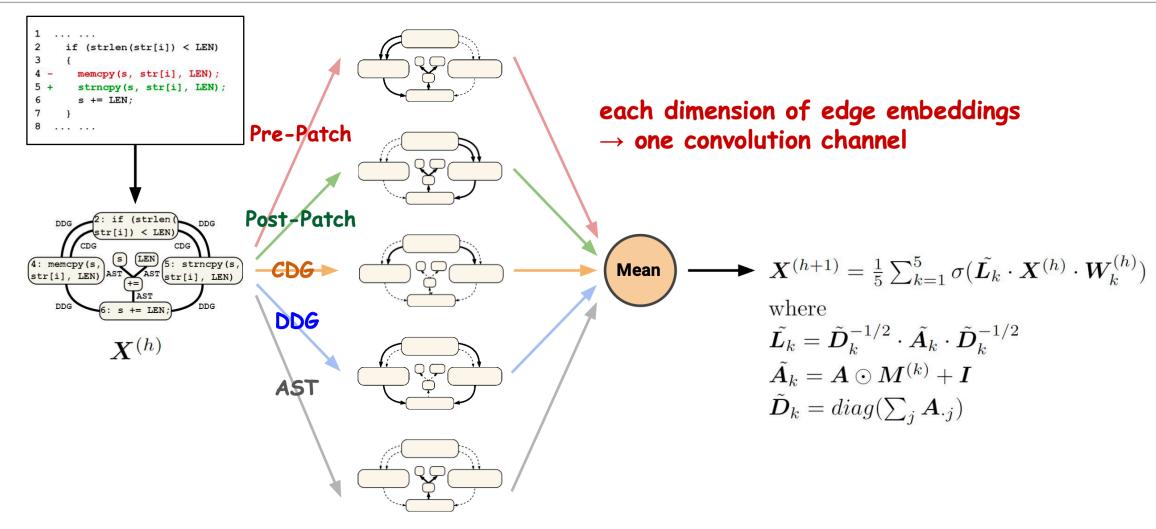
- Edge Embedding
 - o 5-dimensional binary vector.



e.g., [1,1,0,1,0] means the edge is a context edge of data dependency.

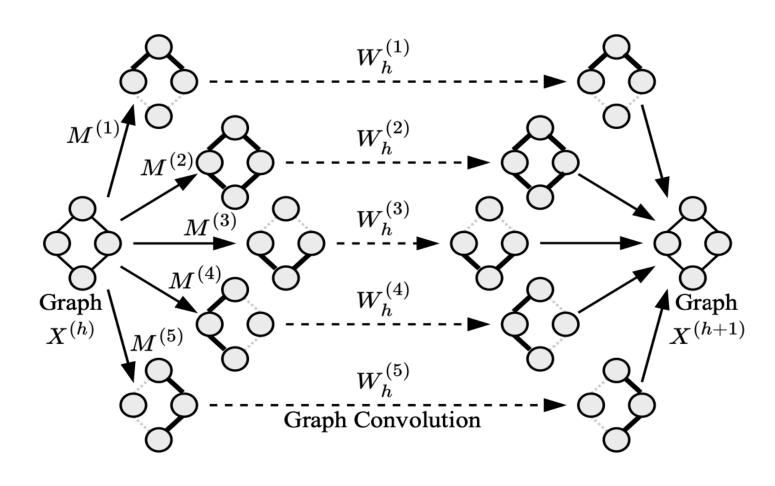
PatchGNN with Multi-Attribute Graph Convolution





PatchGNN with Multi-Attribute Graph Convolution





Implementation & Evaluation



Implementation

• **5K new LoC** in Scala and Python on top of *Joern* parser and *PyTorch* library.

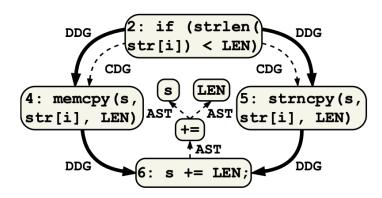
Datasets:

- PatchDB: 12K security patches from 300+ GitHub repos.
- SPI-DB: 10K security patches from FFmpeg and QEMU.

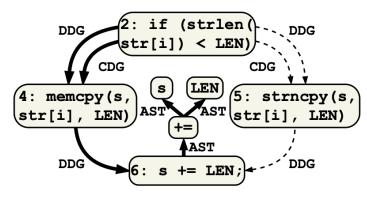
Evaluation:

- Compared with sequential-based patch detector.
- Compared with vulnerability detection methods.
- Case study on four popular OSS repos.

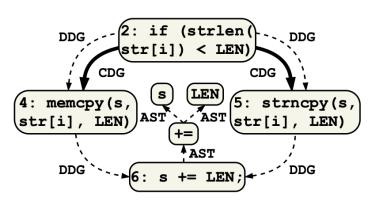
(a) Patch Code.



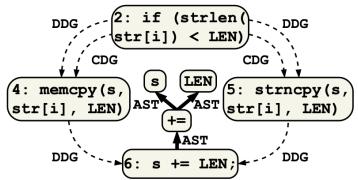
(c) Data Dependency Subgraph.



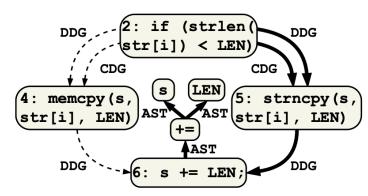
(e) Pre-patch Subgraph.



(b) Control Dependency Subgraph.



(d) Abstract Syntax Tree Subgraph.



(f) Post-patch Subgraph.

[GraphSPD]

Compared with Sequential-based Solution



Accuracy 10.8%↑

F-1 score: 0.096↑

Precision: 28.82%↑

False Positive Rate: 14.62%↓

Method	Dataset	General Metrics		Special Metrics	
Wittild		Accuracy	F1-score	Precision	FP Rate
TwinRNN	PatchDB	69.60%	0.461	48.45%	19.67%
[1][2]	SPI-DB	56.37%	0.512	49.07%	41.57%
GraphSPD	PatchDB	80.39%	0.557	77.27%	5.05%
	SPI-DB	63.04%	0.503	63.96%	19.16%

^[1] PatchRNN: A Deep Learning-Based System for Security Patch Identification.

^[2] SPI: Automated Identification of Security Patches via Commits.

Compared with Vulnerability Detection Solutions



• 2.5 - 50x detection rate of vulnerability detectors.

Method	# Vul _{pre-patch}	# Vul _{post-patch}	# Patch _{security}	TP Rate
Cppcheck[3]	3	1	2	0.54%
flawfinder[4]	109	108	1	0.27%
ReDeBug[5]	29	29	0	0.00%
VUDDY[6]	22	16	21	5.71%
VulDeePecker[7]	3	0	3	0.82%
GraphSPD	-	-	53	14.40%

^[3] Cppcheck. https://cppcheck.sourceforge.io.

^[4] flawfinder. https://dwheeler.com/flawfinder/.

^[5] Redebug: finding unpatched code clones in entire os distributions.

^[6] VUDDY: A scalable approach for vulnerable code clone discovery.

^[7] VulDeePecker: A deep learning- based system for vulnerability detection.

Case Study on False Negetives



Listing 4: Security patch for a double free (CVE-2011-3934).

```
commit 360e95d45ac4123255a4c796db96337f332160ad
if (priv->cac_id_len) {
    serial->len=MIN(priv->cac_id_len, SC_MAX_SERIALNR);
    memcpy(serial->val,priv->cac_id,priv->cac_id_len);
    memcpy(serial->val,priv->cac_id,serial->len);
    SC_RETURN(card->ctx,SC_DEBUG_NORMAL,SC_SUCCESS);
}
```

Listing 5: A patch with similar patterns (CVE-2018-16393).

Case Study on OSS Repos



NGINX: detect 21 security patches (Precision: 78%).

Changes w/	CVE	Total Commits	Valid Commits	Detected S.P.	Confirmed S.P.	Precision
1.19.x 1.17.x 1.15.x 1.13.x	3 3 1 1	180 134 203 270	127 82 120 157	7 4 7 9	6 3 4 8	86% 75% 57% 89%
Sum.	8	787	486	27	21	78%

Xen: detect 29 security patches (Precision: 55%).

OpenSSL: detect 45 security patches (Precision: 66%).

ImageMagick: detect 6 security patches (Precision: 46.2%).

Slicing Depth



Slicing	General	Metrics	Special Metrics		
Iteration No. (N)	Accuracy	F1-score	Precision	FP Rate	
0	80.19%	0.555	76.11%	5.42%	
1	80.39%	0.557	77.27%	5.05%	
2	79.24%	0.531	73.61%	5.87%	
	76.90%	0.501	64.42%	8.95%	

Conclusion



- Silent security patches can be leveraged by attackers to launch N-day attacks.
- GraphSPD presents patches as graphs and identifies security patches with graph learning, achieving higher accuracy and fewer false alarms.
- GraphSPD can be extended to other programming languages.

Discuss



- PatchDB only collects security patches from NVD
- SPI-DB contains security patches only from two repositories
- GraphSPD mainly learns from the existing patterns and may not apply to the unseen ones
- Cannot distinguish specific security patch types:
 - Some patches are too rare to be included in training set
 - Data is imbalanced in security patch datasets. According to the analysis on NVD, 24.6% of vulnerabilities are related to code execution, whereas only 0.1% of vulnerabilities are HTTP response splitting
- Single within function support!

Acknowledgments



- [GraphSPD] GraphSPD: Graph-Based Security Patch Detection with Enriched Code Semantics, S. Wang, X. Wang, K. Sun, S. Jajodia, H. Wang, and Q. Li, IEEE S&P 2023.
- [CPG] Modeling and Discovering Vulnerabilities with Code Property Graphs, F. Yamaguchi, N. Golde, D. Arp, and K. Rieck, IEEE S&P 2014.