



40-414 Compiler Design

Local Optimizations

Lecture 11

Exercise

Question?

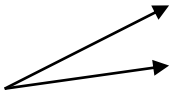
Which of the following are valid local optimizations for the given basic block? Assume that only g and x are referenced outside of this basic block.

- Copy propagation: Line 4 becomes $d := a * b$.
- Common subexpression elimination: Line 5 becomes $e := d$.
- Dead code elimination: Line 3 is removed.
- After many rounds of valid optimizations, the entire block can be reduced to $g := 5$.

```
1  a := 1
2  b := 3
3  c := a + x
4  d := a * 3
5  e := b * 3
6  f := a + b
7  g := e - f
```

Question?

Which of the following are valid local optimizations for the given basic block? Assume that only g and x are referenced outside of this basic block.

Copy propagation 

```
1  a := 1
2  b := 3
3  c := a + x
4  d := a * 3
5  e := b * 3
6  f := a + b
7  g := e - f
```

- Copy propagation: Line 4 becomes $d := a * b$.
- Common subexpression elimination: Line 5 becomes $e := d$.
- Dead code elimination: Line 3 is removed.
- After many rounds of valid optimizations, the entire block can be reduced to $g := 5$.

Question?

Which of the following are valid local optimizations for the given basic block? Assume that only g and x are referenced outside of this basic block.

- Copy propagation: Line 4 becomes $d := a * b$.
- Common subexpression elimination: Line 5 becomes $e := d$.
- Dead code elimination: Line 3 is removed.
- After many rounds of valid optimizations, the entire block can be reduced to $g := 5$.

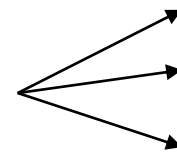
```
1  a := 1
2  b := 3
3  c := 1 + x
4  d := 1 * 3
5  e := 3 * 3
6  f := 1 + 3
7  g := e - f
```

Question?

Which of the following are valid local optimizations for the given basic block? Assume that only g and x are referenced outside of this basic block.

- Copy propagation: Line 4 becomes $d := a * b$.
- Common subexpression elimination: Line 5 becomes $e := d$.
- Dead code elimination: Line 3 is removed.
- After many rounds of valid optimizations, the entire block can be reduced to $g := 5$.

Constant folding



```
1  a := 1
2  b := 3
3  c := 1 + x
4  d := 1 * 3
5  e := 3 * 3
6  f := 1 + 3
7  g := e - f
```

Question?

Which of the following are valid local optimizations for the given basic block? Assume that only g and x are referenced outside of this basic block.

- Copy propagation: Line 4 becomes $d := a * b$.
- Common subexpression elimination: Line 5 becomes $e := d$.
- Dead code elimination: Line 3 is removed.
- After many rounds of valid optimizations, the entire block can be reduced to $g := 5$.

```
1  a := 1
2  b := 3
3  c := 1 + x
4  d := 3
5  e := 9
6  f := 4
7  g := e - f
```

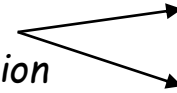
Question?

Which of the following are valid local optimizations for the given basic block? Assume that only g and x are referenced outside of this basic block.

- Copy propagation: Line 4 becomes $d := a * b$.
- Common subexpression elimination: Line 5 becomes $e := d$.
- Dead code elimination: Line 3 is removed.
- After many rounds of valid optimizations, the entire block can be reduced to $g := 5$.

```
1  a := 1
2  b := 3
3  c := 1 + x
4  d := 3
5  e := 9
6  f := 4
7  g := e - f
```

Copy propagation



Question?

Which of the following are valid local optimizations for the given basic block? Assume that only g and x are referenced outside of this basic block.

- Copy propagation: Line 4 becomes $d := a * b$.
- Common subexpression elimination: Line 5 becomes $e := d$.
- Dead code elimination: Line 3 is removed.
- After many rounds of valid optimizations, the entire block can be reduced to $g := 5$.

```
1  a := 1
2  b := 3
3  c := 1 + x
4  d := 3
5  e := 9
6  f := 4
7  g := 9 - 4
```

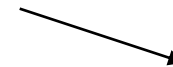

Question?

Which of the following are valid local optimizations for the given basic block? Assume that only g and x are referenced outside of this basic block.

- Copy propagation: Line 4 becomes $d := a * b$.
- Common subexpression elimination:
Line 5 becomes $e := d$.
- Dead code elimination: Line 3 is removed.
- After many rounds of valid optimizations, the entire block can be reduced to $g := 5$.

```
1  a := 1
2  b := 3
3  c := 1 + x
4  d := 3
5  e := 9
6  f := 4
7  g := 9 - 4
```

*Constant
folding*



Question?

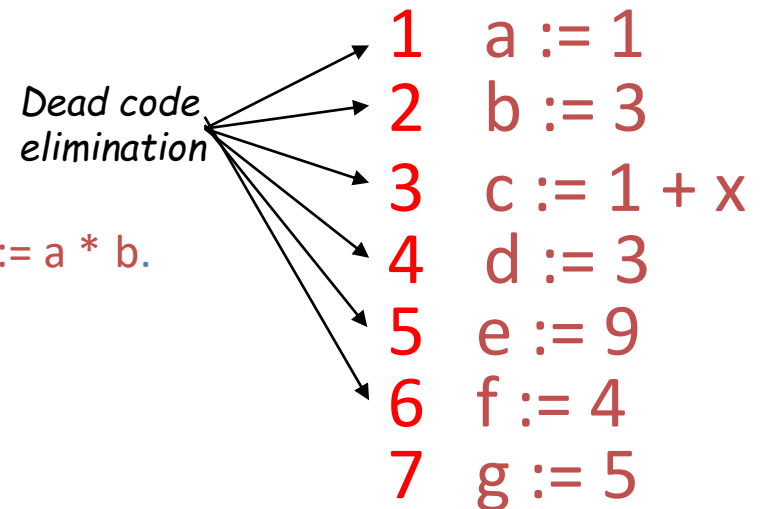
Which of the following are valid local optimizations for the given basic block? Assume that only g and x are referenced outside of this basic block.

- Copy propagation: Line 4 becomes $d := a * b$.
- Common subexpression elimination: Line 5 becomes $e := d$.
- Dead code elimination: Line 3 is removed.
- After many rounds of valid optimizations, the entire block can be reduced to $g := 5$.

```
1  a := 1
2  b := 3
3  c := 1 + x
4  d := 3
5  e := 9
6  f := 4
7  g := 5
```

Question?

Which of the following are valid local optimizations for the given basic block? Assume that only g and x are referenced outside of this basic block.



- Copy propagation: Line 4 becomes $d := a * b$.
- Common subexpression elimination: Line 5 becomes $e := d$.
- Dead code elimination: Line 3 is removed.
- After many rounds of valid optimizations, the entire block can be reduced to $g := 5$.

Question?

Which of the following are valid local optimizations for the given basic block? Assume that only g and x are referenced outside of this basic block.

- Copy propagation: Line 4 becomes $d := a * b$.
- Common subexpression elimination: Line 5 becomes $e := d$.
- Dead code elimination: Line 3 is removed.
- After many rounds of valid optimizations, the entire block can be reduced to $g := 5$.

1
2
3
4
5
6
7 $g := 5$

Answer!

Which of the following are valid local optimizations for the given basic block? Assume that only g and x are referenced outside of this basic block.

- Copy propagation: Line 4 becomes $d := a * b$.
- Common subexpression elimination: Line 5 becomes $e := d$.
- Dead code elimination: Line 3 is removed.
- After many rounds of valid optimizations, the entire block can be reduced to $g := 5$.

```
1  a := 1
2  b := 3
3  c := a + x
4  d := a * 3
5  e := b * 3
6  f := a + b
7  g := e - f
```

Example: C code

```
void quicksort(m, n)
    int m, n;
    {
        int i, j;
        if (n <= m ) return;
        /* fragment begins here */
        i = m-1; j = n; v = a[n];
        while(1)    {
            do i = i+1; while( a[i] < v );
            do j = j-1; while( a[j] > v );
            if( i >= j ) break;
            x = a[i]; a[i] = a[j]; a[j] = x;
        }
        x = a[i]; a[i] = a[n]; a[n]= x;
        /* fragment ends here */
        quicksort(m, j); quicksort(i+1, n);
    }
```

Augmented 3AC

An augmented 3 address code language to simplify the code...

Let a be an array of integers starting at byte address a_0

$a[\text{add}]$ on the left-hand-side of an assignment is the address $a_0 + \text{add}$

$a[\text{add}]$ on the right-hand-side of an assignment is the value of the element of the array at address $a_0 + \text{add}$

*Since integers are stored in 4 bytes the offset address of an element $a[i]$ is $4 * i$*

Augmented 3AC of the C code

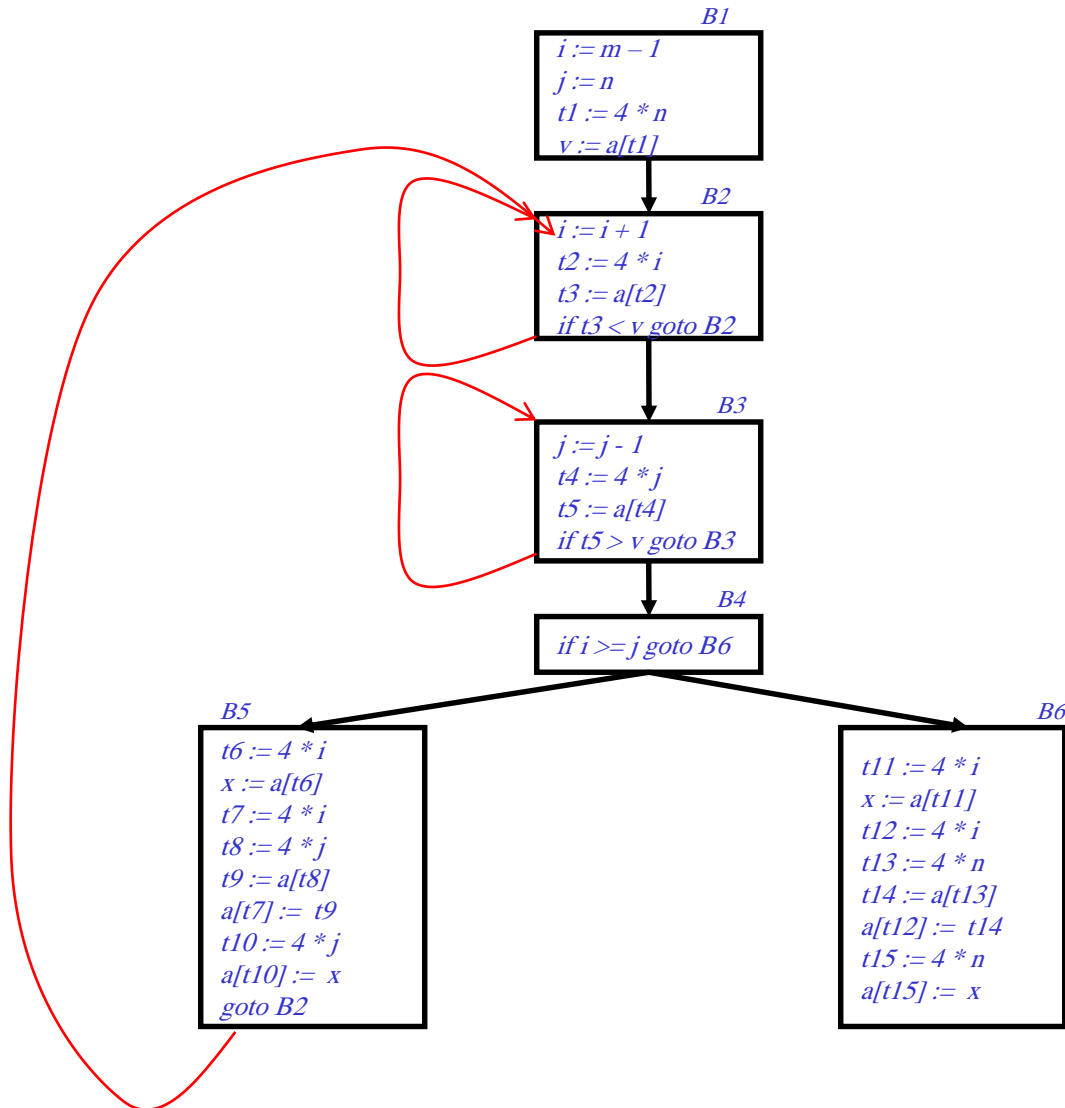
```
01)  $i := m - 1$ 
02)  $j := n$ 
03)  $t1 := 4 * n$ 
04)  $v := a[t1]$ 
05)  $i := i + 1$ 
06)  $t2 := 4 * i$ 
07)  $t3 := a[t2]$ 
08) if  $t3 < v$  goto 5
09)  $j := j - 1$ 
10)  $t4 := 4 * j$ 
11)  $t5 := a[t4]$ 
12) if  $t5 > v$  goto 9
13) if  $i \geq j$  goto 23
14)  $t6 := 4 * i$ 
15)  $x := a[t6]$ 
16)  $t7 := 4 * i$ 
17)  $t8 := 4 * j$ 
18)  $t9 := a[t8]$ 
19)  $a[t7] := t9$ 
20)  $t10 := 4 * j$ 
21)  $a[t10] := x$ 
22) goto 5
23)  $t11 := 4 * i$ 
24)  $x := a[t11]$ 
25)  $t12 := 4 * i$ 
26)  $t13 := 4 * n$ 
27)  $t14 := a[t13]$ 
28)  $a[t12] := t14$ 
29)  $t15 := 4 * n$ 
30)  $a[t15] := x$ 
```


Basic Blocks

01) $i := m - 1$
02) $j := n$
03) $t1 := 4 * n$
04) $v := a[t1]$
05) $i := i + 1$
06) $t2 := 4 * i$
07) $t3 := a[t2]$
08) if $t3 < v$ goto 5
09) $j := j - 1$
10) $t4 := 4 * j$
11) $t5 := a[t4]$
12) if $t5 > v$ goto 9
13) if $i \geq j$ goto 23
14) $t6 := 4 * i$
15) $x := a[t6]$

16) $t7 := 4 * i$
17) $t8 := 4 * j$
18) $t9 := a[t8]$
19) $a[t7] := t9$
20) $t10 := 4 * j$
21) $a[t10] := x$
22) goto 5
23) $t11 := 4 * i$
24) $x := a[t11]$
25) $t12 := 4 * i$
26) $t13 := 4 * n$
27) $t14 := a[t13]$
28) $a[t12] := t14$
29) $t15 := 4 * n$
30) $a[t15] := x$

Control Flow Graph



Local Optimizations

B5 before

```
t6 := 4 * i
x := a[t6]
t7 := 4 * i
t8 := 4 * j
t9 := a[t8]
a[t7] := t9
t10 := 4 * j
a[t10] := x
goto B2
```

B5 after

```
t6 := 4 * i
x := a[t6]
t7 := t6
t8 := 4 * j
t9 := a[t8]
a[t7] := t9
t10 := t8
a[t10] := x
goto B2
```

Common Subexpression Elimination

Local Optimizations

B5 before

```
t6 := 4 * i
x := a[t6]
t7 := t6
t8 := 4 * j
t9 := a[t8]
a[t7] := t9
t10 := t8
a[t10] := x
goto B2
```

B5 after

```
t6 := 4 * i
x := a[t6]
t7 := t6
t8 := 4 * j
t9 := a[t8]
a[t6] := t9
t10 := t8
a[t8] := x
goto B2
```

Copy propagation

Local Optimizations

B5 before

```
t6 := 4 * i
x := a[t6]
t7 := t6
t8 := 4 * j
t9 := a[t8]
a[t7] := t9
t10 := t8
a[t10] := x
goto B2
```

B5 after

```
t6 := 4 * i
x := a[t6]
t8 := 4 * j
t9 := a[t8]
a[t6] := t9
a[t8] := x
goto B2
```

Dead code elimination

Local Optimizations

B6 before

```
t11 := 4 * i
x := a[t11]
t12 := 4 * i
t13 := 4 * n
t14 := a[t13]
a[t12] := t14
t15 := 4 * n
a[t15] := x
```

B6 after

```
t11 := 4 * i
x := a[t11]
t12 := t11
t13 := 4 * n
t14 := a[t13]
a[t12] := t14
t15 := t13
a[t15] := x
```

Common Subexpression Elimination

Local Optimizations

B6 before

```
t11 := 4 * i
x := a[t11]
t12 := t11
t13 := 4 * n
t14 := a[t13]
a[t12] := t14
t15 := t13
a[t15] := x
```

B6 after

```
t11 := 4 * i
x := a[t11]
t12 := t11
t13 := 4 * n
t14 := a[t13]
a[t11] := t14
t15 := t13
a[t13] := x
```

Copy Propagation

Local Optimizations

B6 before

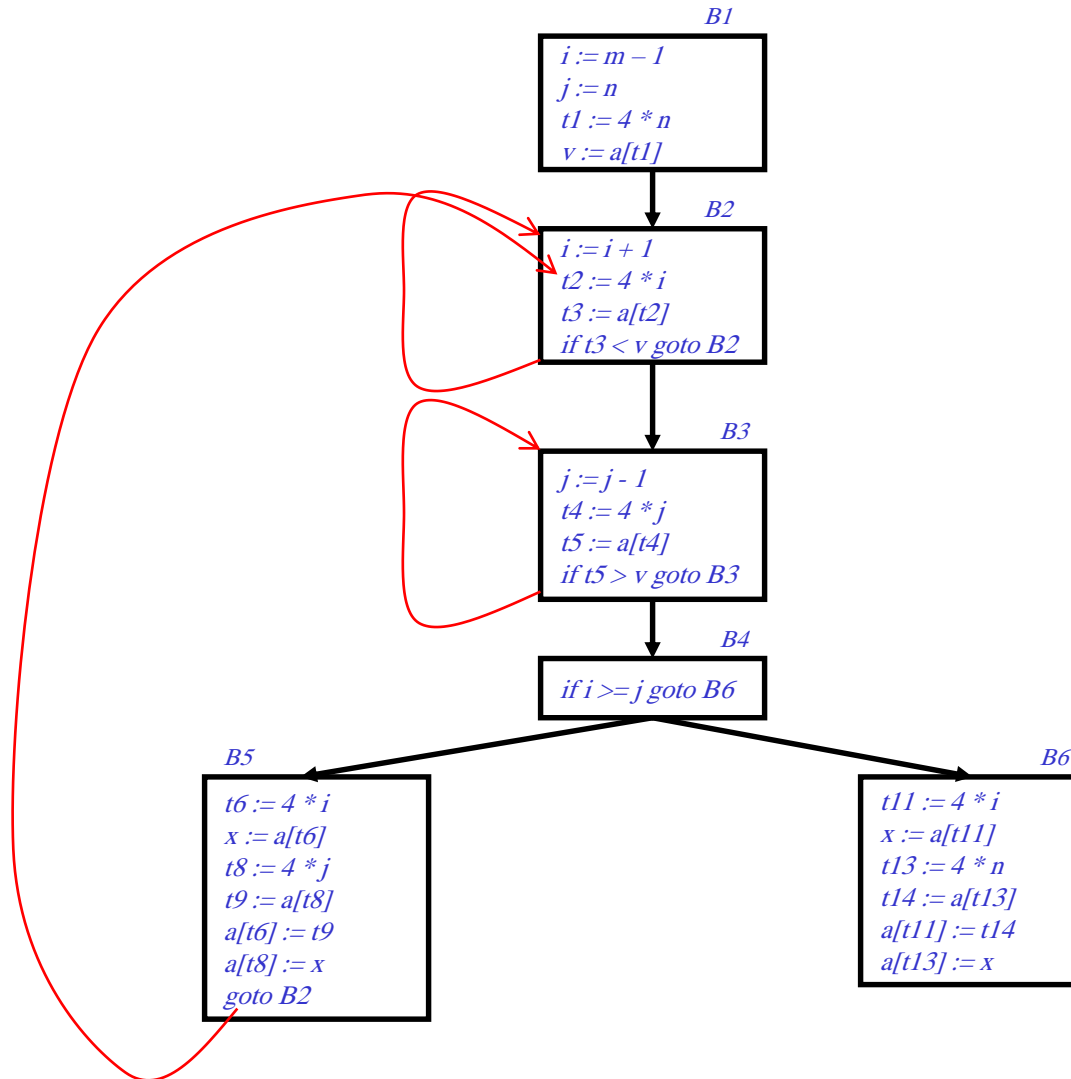
```
t11 := 4 * i
x := a[t11]
t12 := t11
t13 := 4 * n
t14 := a[t13]
a[t11] := t14
t15 := t13
a[t13] := x
```

Dead code elimination

B6 after

```
t11 := 4 * i
x := a[t11]
t13 := 4 * n
t14 := a[t13]
a[t11] := t14
a[t13] := x
```


After Local Optimizations



Reduction in Strength

In *B2* whenever i increases by 1, $t2$ increases by 4

In *B3* whenever j decreases by 1, $t4$ decreases by 4

B1 Before

```
 $i := m - 1$   
 $j := n$   
 $t1 := 4 * n$   
 $v := a[t1]$ 
```

B2:

```
 $i := i + 1$   
 $t2 := 4 * i$   
 $t3 := a[t2]$   
if  $t3 < v$  goto B2
```

B3:

```
 $j := j - 1$   
 $t4 := 4 * j$   
 $t5 := a[t4]$   
if  $t5 > v$  goto B2
```

B1 After

```
 $i := m - 1$   
 $j := n$   
 $t1 := 4 * n$   
 $v := a[t1]$   
 $t2 := 4 * i$   
 $t4 := 4 * j$ 
```

B2:

```
 $i := i + 1$   
 $t2 := t2 + 4$   
 $t3 := a[t2]$   
if  $t3 < v$  goto B2
```

B3:

```
 $j := j - 1$   
 $t4 := t4 - 4$   
 $t5 := a[t4]$   
if  $t5 > v$  goto B3
```

Induction Variables Elimination

In *B2* whenever i increases by 1, $t2$ increases by 4,
 i and $t2$ are called induction variables.

In *B3* whenever j decreases by 1, $t4$ decreases by 4,
 j and $t4$ are induction variables, too.

If there are two or more induction variables in a loop, it
may be possible to get rid of all but one

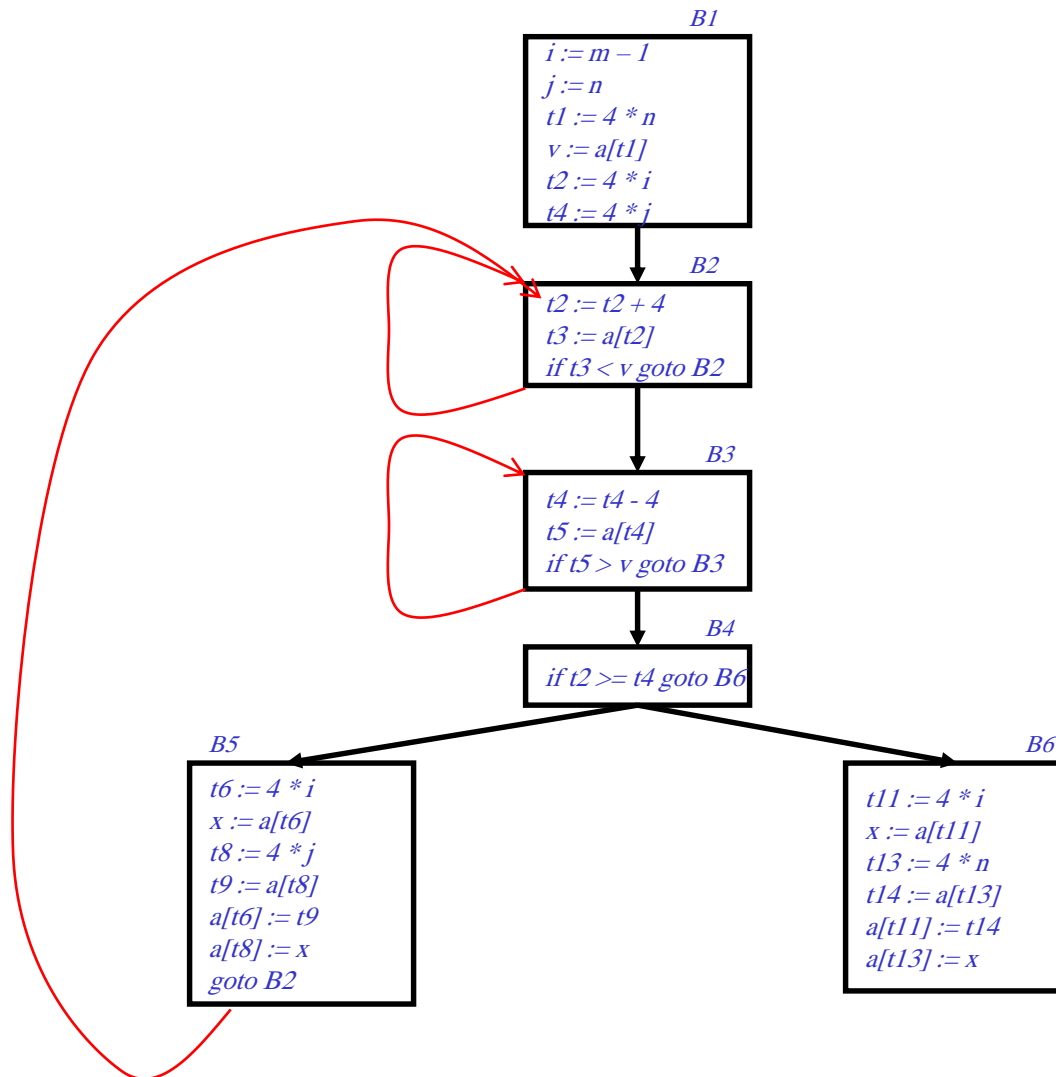
B4: Before

if $i \geq j$ goto B6

B4: After

if $t2 \geq t4$ goto B6

After Loop Optimizations



After Global Optimizations

