



40-414 Compiler Design

Introduction to Parsing

Lecture 4

Exercise

Question?

- Consider the grammar

$$E \rightarrow TX$$

$$T \rightarrow (E) \mid \text{int } Y$$

$$X \rightarrow + E \mid \varepsilon$$

$$Y \rightarrow * T \mid \varepsilon$$

- Write Recursive Descent Procedures including panic mode error recovery for all non-terminals.
- Write a step-by-step parsing of input 'int * int'
- Draw the parse tree of the input

Procedure E

Follow(E) = {), \$ }

- Consider the grammar

$E \rightarrow TX$

$X \rightarrow + E \mid \varepsilon$

$T \rightarrow (E) \mid \text{int } Y$

$Y \rightarrow * T \mid \varepsilon$

procedure E ;

{ if lookahead is in { (, int }

 then { call T; call X }

 else if lookahead is in { \$,) }

 then { print ('missing E on line ...'); exit }

 else { print ('illegal lookahead on line ...');

 lookahead := get_next_token;

 call E

 }

}

Procedure T

Follow(**T**) = { +,), \$ }

- Consider the grammar

$E \rightarrow TX$

$X \rightarrow + E \mid \varepsilon$

$T \rightarrow (E) \mid \text{int } Y$

$Y \rightarrow * T \mid \varepsilon$

procedure **T** ;

{ if lookahead = '('

 then { call Match ('('); call **E**; call Match(')'); }

 else if lookahead = **int**

 then { call Match (**int**); call **Y** }

 else if lookahead is in { + , \$,) }

 then { print ('missing **T** on line ...'); **exit** }

 else { print ('illegal lookahead on line ...');

 lookahead := get_next_token;

 call **T** }

}

Procedure X

Follow(X) = {), \$ }

- Consider the grammar

$E \rightarrow TX$

$T \rightarrow (E) \mid \text{int } Y$

$X \rightarrow + E \mid \varepsilon$

$Y \rightarrow * T \mid \varepsilon$

procedure X ;

{ if lookahead = '+'

then { call Match ('+'); call E }

else if lookahead is in { \$,) }

then exit;

else { print ('illegal lookahead on line ...');

lookahead := get_next_token;

call X }

}

Procedure Y

Follow(Y) = { +,), \$ }

- Consider the grammar

$E \rightarrow TX$

$T \rightarrow (E) \mid \text{int } Y$

$X \rightarrow + E \mid \varepsilon$

$Y \rightarrow * T \mid \varepsilon$

procedure Y ;

{ if *lookahead* = '*'

then { call Match ('*'); call T }

else if *lookahead* is in { \$,), + }

then exit;

else { print ('illegal *lookahead* on line ...');

lookahead := get_next_token;

call Y }

}

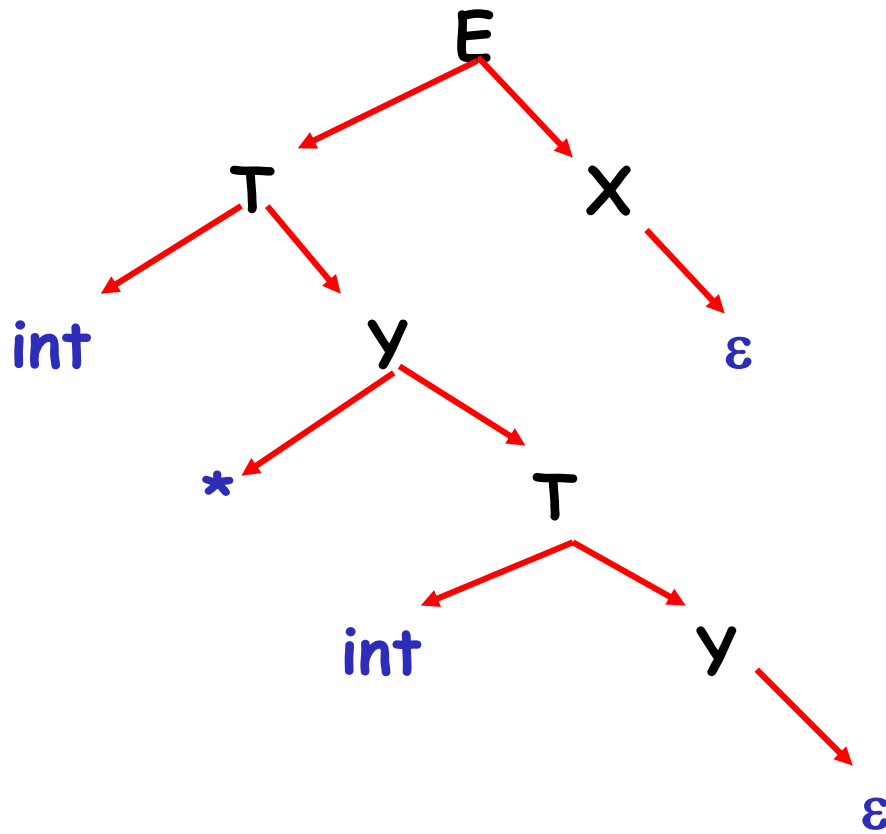
R.D. Parsing Example

$$\begin{aligned}
 E &\rightarrow TX \\
 T &\rightarrow (E) \mid \text{int } Y \\
 X &\rightarrow + E \mid \varepsilon \\
 Y &\rightarrow * T \mid \varepsilon
 \end{aligned}$$

Current Routine		Input	Action
E	→	int * int \$	call T
T	→	int * int \$	Match (int)
Match	→	int * int \$	Call Lexer
T	←	* int \$	call Y
Y	→	* int \$	Match (*)
Match	→	int \$	Call Lexer
Y	←	int \$	call T
T	→	int \$	Match (int)
Match	→	int \$	Call Lexer
T	←	\$	call Y
Y	→	\$	exit
T	←	\$	exit
Y	←	\$	exit
T	←	\$	exit
E	←	\$	exit
X	→	\$	exit
E	←	\$	ACCEPT

Calls: →
 Returns: ←

Parsing Tree of $\text{int} * \text{int} \$$



Current Routine

E	→
T	→
Match	→
T	←
Y	→
Match	→
Y	←
T	→
Match	→
T	←
Y	→
T	←
Y	←
T	←
E	←
X	→
E	←

Question?

How many strings does the following grammar generate?

- 7
- 15
- 2
- 8
- 16
- 4

$$\begin{aligned}A &\rightarrow B B \\B &\rightarrow C C \\C &\rightarrow 1 \mid 2\end{aligned}$$

Answer!

How many strings does the following grammar generate?

7

15

2

8

16

4

$A \rightarrow B B$

$B \rightarrow C C$

$C \rightarrow 1 \mid 2$

Question?

How many strings does the following grammar generate?

- 16
- 31
- 15
- 12
- 64
- 63
- 32
- 11

$$A \rightarrow B B$$

$$B \rightarrow C C$$

$$C \rightarrow 1 \mid 2 \mid \varepsilon$$

Answer!

How many strings does the following grammar generate?

$$A \rightarrow B B$$

$$B \rightarrow C C$$

$$C \rightarrow 1 \mid 2 \mid \varepsilon$$

16

31

15

12

64

63

32

11

Question?

Which of the following is a valid derivation of the given grammar?

$$S \rightarrow aXa$$

$$X \rightarrow \varepsilon \mid bY$$

$$Y \rightarrow \varepsilon \mid cXc \mid d$$

S
aXa
abYa
acXca
acca

S
aa

S
aXa
abYa
abcXca
abcbYca
abcbdca

S
aXa
abYa
abcXcda
abccda

Answer!

Which of the following is a valid derivation of the given grammar?

S
aXa
abYa
acXca
acca

S
aa

S
aXa
abYa
abcXca
abcbYca
abcbdca

S
aXa
abYa
abcXcda
abccda

$S \rightarrow aXa$

$X \rightarrow \varepsilon \mid bY$

$Y \rightarrow \varepsilon \mid cXc \mid d$

Derivation:

$S \rightarrow aXa \rightarrow abYa$

$\rightarrow abcXca \rightarrow abcbYca$

$\rightarrow abcbdca$

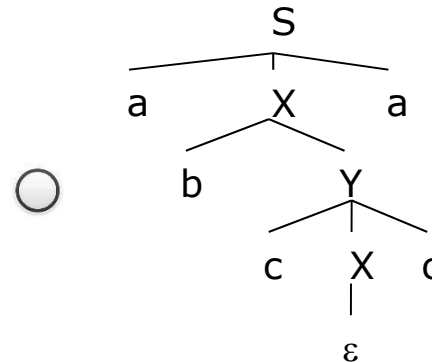
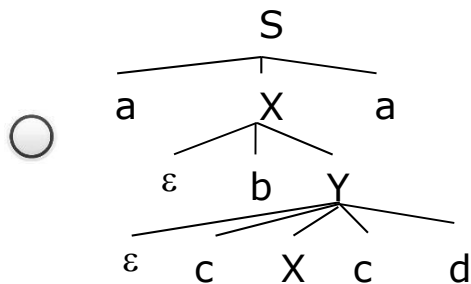
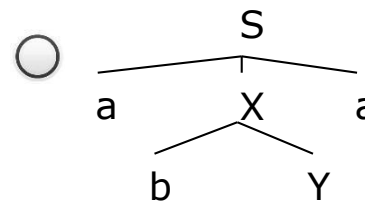
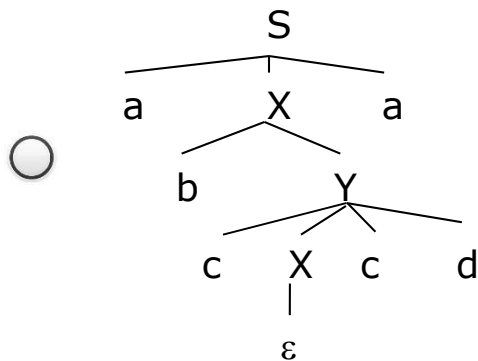
Question?

Which of the following is a valid parse tree for the given grammar?

$$S \rightarrow aXa$$

$$X \rightarrow \varepsilon \mid bY$$

$$Y \rightarrow \varepsilon \mid cXc \mid d$$



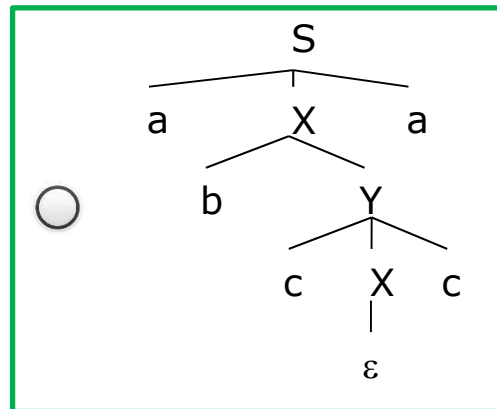
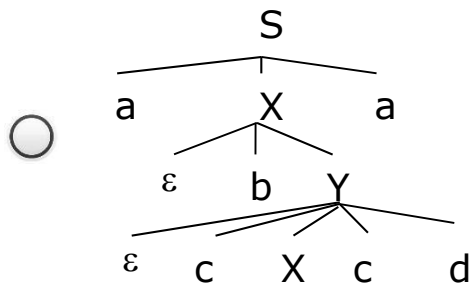
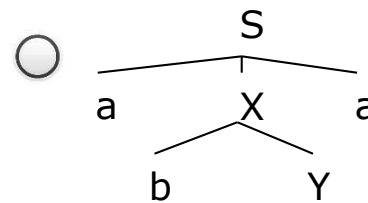
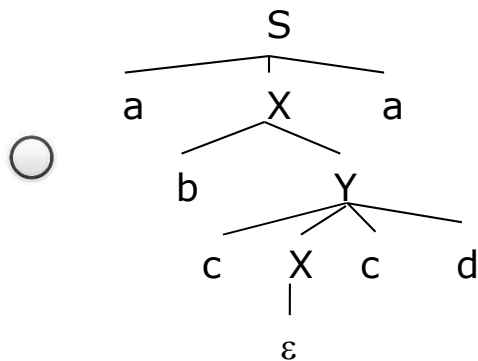
Answer!

Which of the following is a valid parse tree for the given grammar?

$$S \rightarrow aXa$$

$$X \rightarrow \varepsilon \mid bY$$

$$Y \rightarrow \varepsilon \mid cXc \mid d$$



Question?

Choose the grammar that correctly eliminates left recursion from the given grammar:

$$E \rightarrow E + T \mid T$$
$$T \rightarrow id \mid (E)$$

$E \rightarrow E + id \mid E + (E)$
 $\quad \mid id \mid (E)$

$E \rightarrow TE'$
 $E' \rightarrow + TE' \mid \varepsilon$
 $T \rightarrow id \mid (E)$

$E \rightarrow E' + T \mid T$
 $E' \rightarrow id \mid (E)$
 $T \rightarrow id \mid (E)$

$E \rightarrow id + E \mid E + T \mid T$
 $T \rightarrow id \mid (E)$

Answer!

Choose the grammar that correctly eliminates left recursion from the given grammar: $E \rightarrow E + T \mid T$
 $T \rightarrow id \mid (E)$

$E \rightarrow E + id \mid E + (E)$
 $\quad \mid id \mid (E)$

$E \rightarrow TE'$
 $E' \rightarrow + TE' \mid \varepsilon$
 $T \rightarrow id \mid (E)$

$E \rightarrow E' + T \mid T$
 $E' \rightarrow id \mid (E)$
 $T \rightarrow id \mid (E)$

$E \rightarrow id + E \mid E + T \mid T$
 $T \rightarrow id \mid (E)$

Question?

Consider the following grammar. Adding which one of the following rules will cause the grammar to be left-recursive?
[Choose all that apply]

$$S \rightarrow A$$

$$A \rightarrow B \mid C$$

$$B \rightarrow (C)$$

$$C \rightarrow B + C \mid D$$

$$D \rightarrow 1 \mid 0$$

$D \rightarrow A$

$A \rightarrow D$

$B \rightarrow C$

$D \rightarrow B$

$C \rightarrow 1 C$

Answer!

Consider the following grammar. Adding which one of the following rules will cause the grammar to be left-recursive?
[Choose all that apply]

$$S \rightarrow A$$

$$A \rightarrow B \mid C$$

$$B \rightarrow (C)$$

$$C \rightarrow B + C \mid D$$

$$D \rightarrow 1 \mid 0$$

$D \rightarrow A$

$A \rightarrow D$

$B \rightarrow C$

$D \rightarrow B$

$C \rightarrow 1 C$

Question?

Which of the following grammars are ambiguous?

- $S \rightarrow SS \mid a \mid b$
- $E \rightarrow E + E \mid id$
- $S \rightarrow Sa \mid Sb$
- $E \rightarrow E' \mid E' + E$
 $E' \rightarrow -E' \mid id \mid (E)$

Answer!

Which of the following grammars are ambiguous?

$S \rightarrow SS \mid a \mid b$

$E \rightarrow E + E \mid id$

$S \rightarrow Sa \mid Sb$

$E \rightarrow E' \mid E' + E$

$$E' \rightarrow -E' \mid id \mid (E)$$

Question?

Choose the unambiguous version
of the given ambiguous grammar: $S \rightarrow SS | a | b | \varepsilon$

$S \rightarrow Sa | Sb | \varepsilon$

$S \rightarrow SS'$
 $S' \rightarrow a | b$

$S \rightarrow S | S'$
 $S' \rightarrow a | b$

$S \rightarrow Sa | Sb$

Answer!

Choose the unambiguous version
of the given ambiguous grammar: $S \rightarrow SS | a | b | \varepsilon$

$S \rightarrow Sa | Sb | \varepsilon$

$S \rightarrow SS'$
 $S' \rightarrow a | b$

$S \rightarrow S | S'$
 $S' \rightarrow a | b$

$S \rightarrow Sa | Sb$

Question?

Consider the following grammar. How many unique parse trees are there for the string $5 * 3 + (2 * 7) + 4$?

- 2
- 1
- 7
- 8
- 5
- 4

$$E \rightarrow E * E \mid E + E \mid (E) \mid \text{int}$$

Answer!

Consider the following grammar. How many unique parse trees are there for the string $5 * 3 + (2 * 7) + 4$?

2

1

7

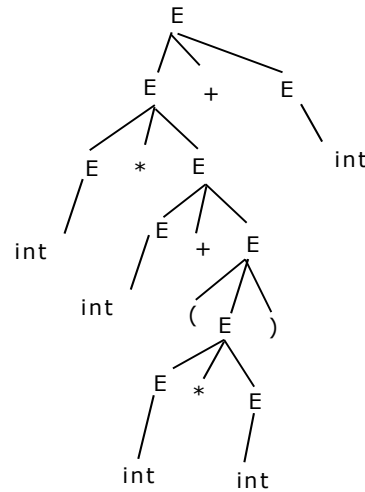
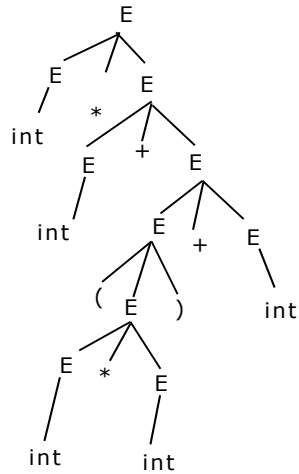
8

5

4

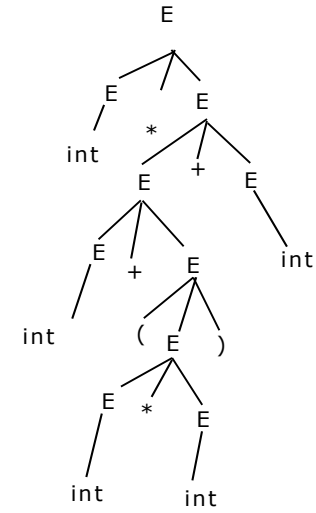
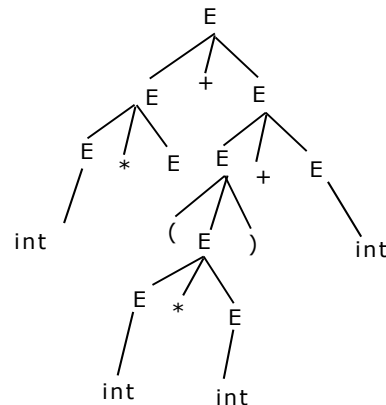
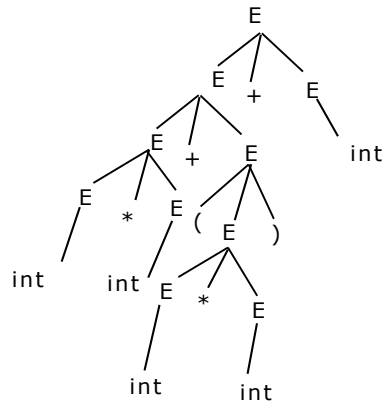
$E \rightarrow E * E \mid E + E \mid (E) \mid \text{int}$

Answer!



$$E \rightarrow E * E \mid E + E \mid (E) \mid int$$

$$5 * 3 + (2 * 7) + 4$$



Question?

Which of the following statements are true about the given grammar?

$$S \rightarrow a T U b \mid \varepsilon$$

$$T \rightarrow c U c \mid b U b \mid a U a$$

$$U \rightarrow S b \mid c c$$

Choose all that are correct.

- The follow set of S is $\{\$, b\}$
- The first set of U is $\{a, b, c\}$
- The first set of S is $\{\varepsilon, a, b\}$
- The follow set of T is $\{a, b, c\}$

Answer!

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$$T \rightarrow c U c \mid b U b \mid a U a$$

$$U \rightarrow S b \mid c c$$

Choose all that are correct.

The follow set of S is $\{\$, b\}$

The first set of U is $\{a, b, c\}$

The first set of S is $\{\varepsilon, a, b\}$

The follow set of T is $\{a, b, c\}$

Question?

Consider the following grammar:

$$\begin{aligned} S &\rightarrow A (S) B \mid \varepsilon \\ A &\rightarrow S \mid S B x \mid \varepsilon \\ B &\rightarrow S B \mid y \end{aligned}$$

What are the first and follow sets of S

- First: $\{x, y, (, \varepsilon\}$ Follow: $\{y, x, (,)\}$
- First: $\{x, \varepsilon\}$ Follow: $\{\$, y, x, (,)\}$
- First: $\{y, (, \varepsilon\}$ Follow: $\{\$, y, (,)\}$
- First: $\{x, y, (, \varepsilon\}$ Follow: $\{\$, y, x, (,)\}$
- First: $\{x, y, (\}$ Follow: $\{\$, y, x, (,)\}$
- First: $\{x, (\}$ Follow: $\{\$, y, x\}$

Answer!

Consider the following grammar:

$$\begin{aligned} S &\rightarrow A (S) B \mid \varepsilon \\ A &\rightarrow S \mid S B x \mid \varepsilon \\ B &\rightarrow S B \mid y \end{aligned}$$

What are the first and follow sets of S

- First: { x, y, (, ε } Follow: { y, x, (,) }
- First: { x, ε } Follow: { \$, y, x, (,) }
- First: { y, (, ε } Follow: { \$, y, (,) }
- First: { x, y, (, ε } Follow: { \$, y, x, (,) }
- First: { x, y, (} Follow: { \$, y, x, (,) }
- First: { x, (} Follow: { \$, y, x }

Question?

Choose the derivation that is a valid recursive descent parse for the string `id + id` in the given grammar. Moves that are followed by backtracking are given in red.

E
E'
E' + E
id + E
id + E'
id + id

E
E' + E
id + E
id + E'
id + id

E
E'
-E'
id
(E)
E' + E
-E' + E
id + E
id + E'
id +-E'
id + id

E
E'
id
E' + E
id + E
id + E'
id + id

←

$E \rightarrow E' \mid E' + E$

$E' \rightarrow -E' \mid id \mid (E)$

Answer!

Choose the derivation that is a valid recursive descent parse for the string `id + id` in the given grammar. Moves that are followed by backtracking are given in red.

- E
 E'
 $E' + E$
 $id + E$
 $id + E'$
 $id + id$

- E
 $E' + E$
 $id + E$
 $id + E'$
 $id + id$

- E
 E'
 $-E'$
 id
 (E)
 $E' + E$
 $-E' + E$
 $id + E$
 $id + E'$
 $id +-E'$
 $id + id$

$$E \rightarrow E' \mid E' + E$$

$$E' \rightarrow -E' \mid id \mid (E)$$

- E
 E'
 id
 $E' + E$
 $id + E$
 $id + E'$
 $id + id$



Question?

Choose the alternative that correctly left factors “if” statements in the given grammar

$EXPR \rightarrow$ if BOOL then { EXPR }
 | if BOOL then { EXPR } else { EXPR }
 | ...
 $BOOL \rightarrow$ true | false

$EXPR \rightarrow$ if true then { EXPR }
 | if false then { EXPR }
 | if true then { EXPR } else { EXPR }
 | if false then { EXPR } else { EXPR }
 | ...

$EXPR \rightarrow$ EXPR' | EXPR' else { EXPR }
 $EXPR' \rightarrow$ if BOOL then { EXPR }
 | ...
 $BOOL \rightarrow$ true | false

$EXPR \rightarrow$ if BOOLEXPR'
 | ...
 $EXPR' \rightarrow$ then { EXPR }
 | then { EXPR } else { EXPR }
 $BOOL \rightarrow$ true | false

$EXPR \rightarrow$ if BOOL then { EXPR } EXPR'
 | ...
 $EXPR' \rightarrow$ else { EXPR } | ϵ
 $BOOL \rightarrow$ true | false

Answer!

Choose the alternative that correctly left factors “if” statements in the given grammar

$EXPR \rightarrow$ if BOOL then { EXPR }
 | if BOOL then { EXPR } else { EXPR }
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 $BOOL \rightarrow$ true | false

$EXPR \rightarrow$ if true then { EXPR }
 | if false then { EXPR }
 | if true then { EXPR } else { EXPR }
 | if false then { EXPR } else { EXPR }
 | ...

$EXPR \rightarrow$ EXPR' | EXPR' else { EXPR }
 $EXPR' \rightarrow$ if BOOL then { EXPR }
 | ...
 $BOOL \rightarrow$ true | false

$EXPR \rightarrow$ if BOOL EXPR'
 | ...
 $EXPR' \rightarrow$ then { EXPR }
 | then { EXPR } else { EXPR }
 $BOOL \rightarrow$ true | false

$EXPR \rightarrow$ if BOOL then { EXPR } EXPR'
 | ...
 $EXPR' \rightarrow$ else { EXPR } | ϵ
 $BOOL \rightarrow$ true | false