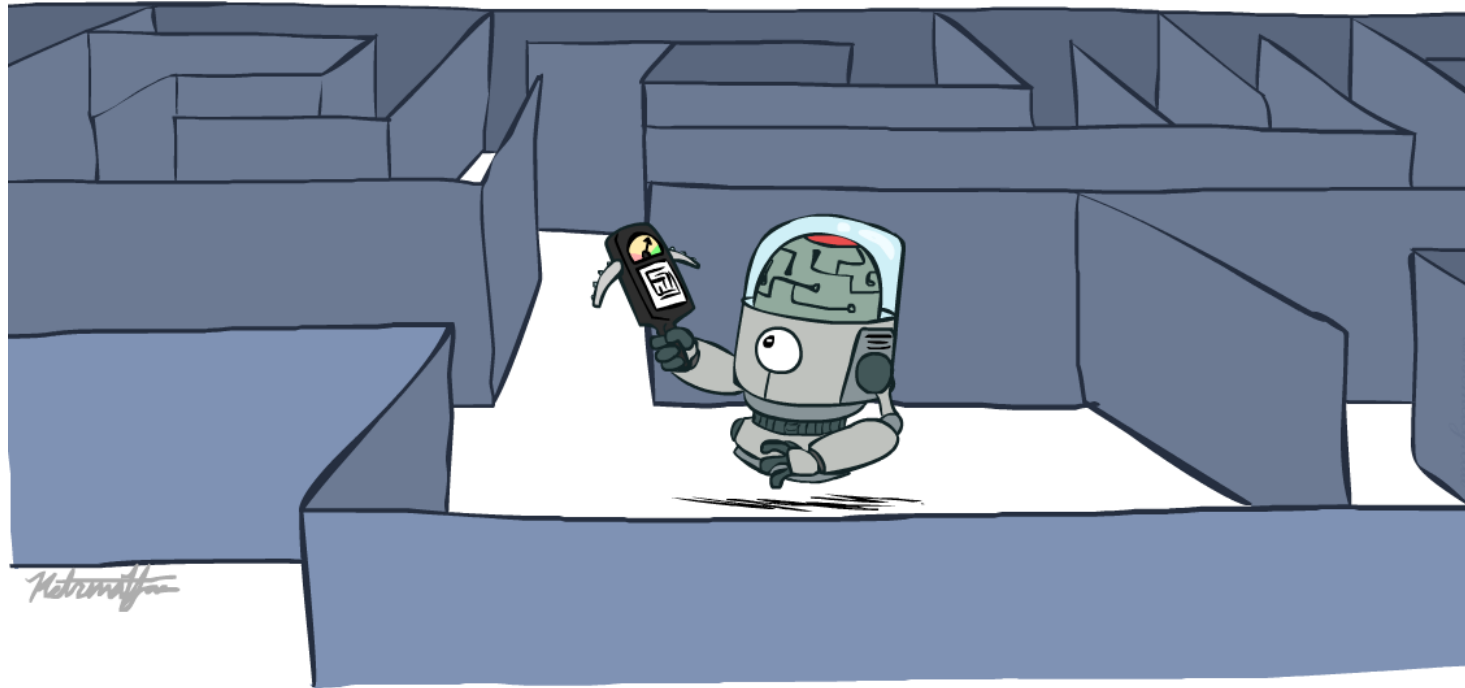


# 40417: Artificial Intelligence

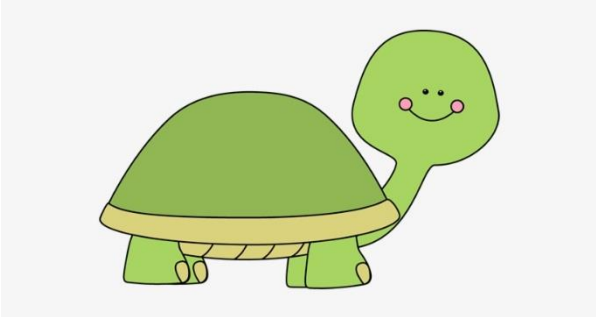
## Informed Search (Part-II)



Instructor: Gholamreza Ghassem-Sani

Sharif University of Technology

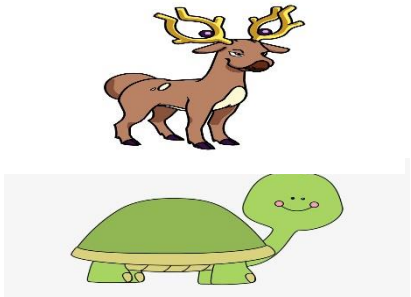
# A\* Search



UCS



Greedy



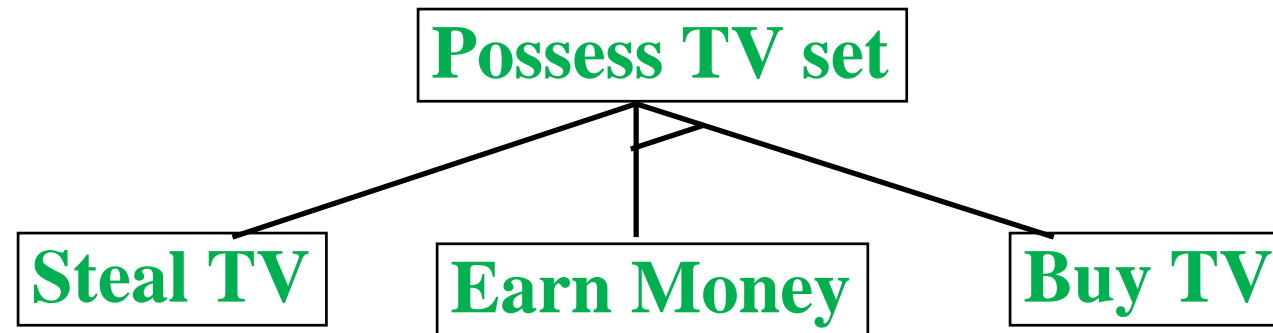
A\*

# Graph Search Pseudo-Code

```
function GRAPH-SEARCH(problem, fringe) return a solution, or failure
  closed ← an empty set
  fringe ← INSERT(MAKE-NODE(INITIAL-STATE[problem]), fringe)
  loop do
    if fringe is empty then return failure
    node ← REMOVE-FRONT(fringe)
    if GOAL-TEST(problem, STATE[node]) then return node
    if STATE[node] is not in closed then
      add STATE[node] to closed
      for child-node in EXPAND(STATE[node], problem) do
        fringe ← INSERT(child-node, fringe)
      end
    end
  end
```

# AND/OR Trees

- Some problems are best represented as achieving subgoals, some of which achieved simultaneously and independently (AND nodes)
- Up to now, we only dealt with OR nodes

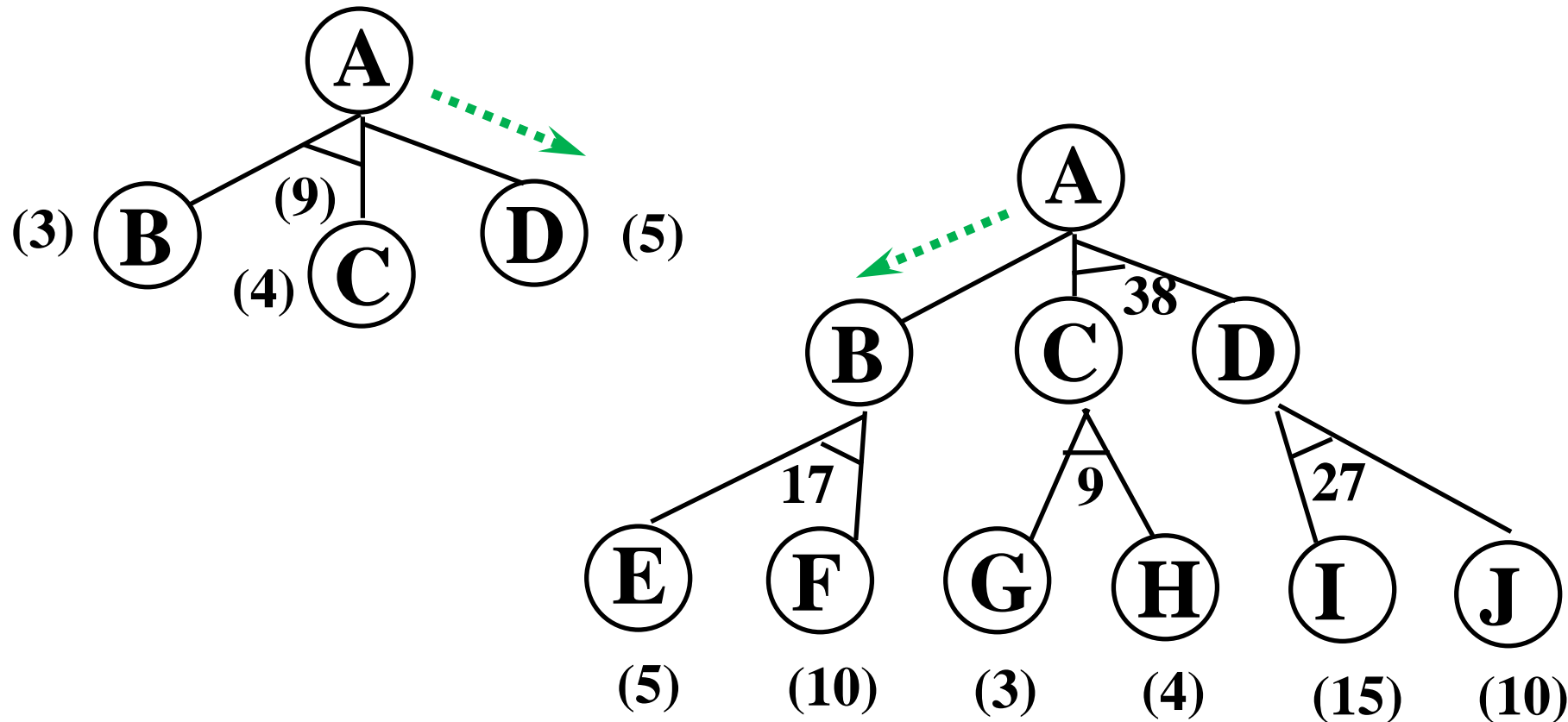


# Searching AND/OR Trees

- A solution in an AND-OR tree is a *sub tree* whose *leafs* are included in the goal set
- Cost function: sum of costs in AND node
$$f(n) = f(n_1) + f(n_2) + \dots + f(n_k)$$
- How can we extend A\* to search AND/OR trees?
  - The AO\* algorithm.

# AND/OR Tree Search

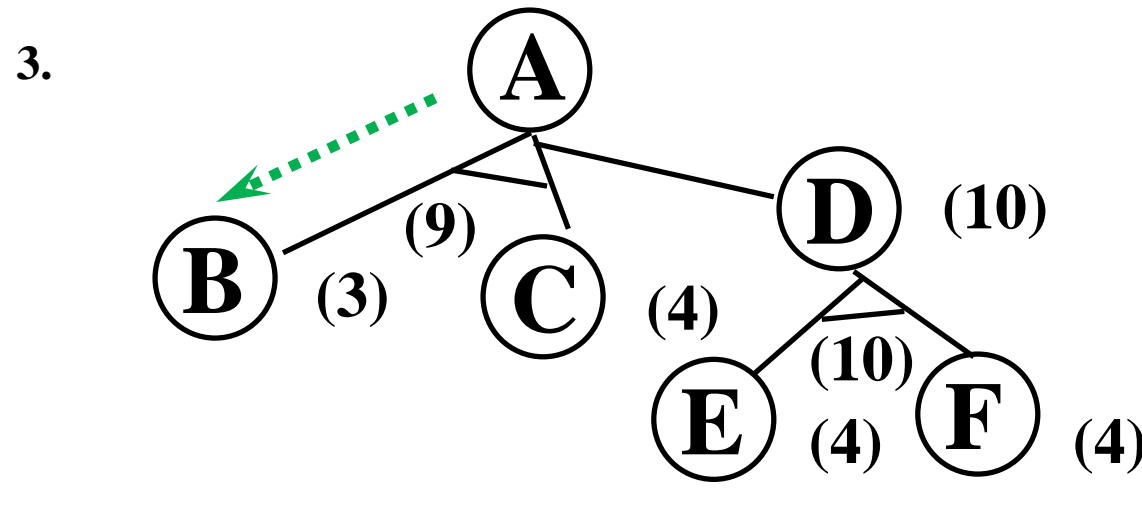
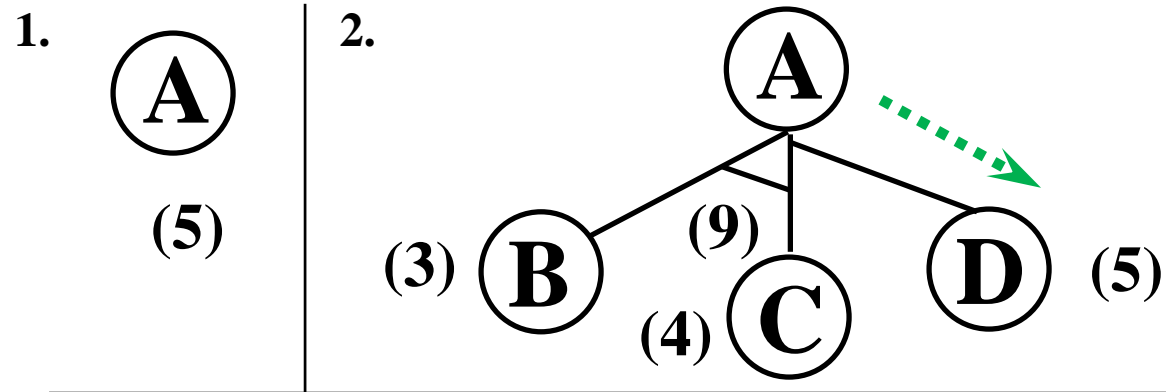
- We must examine several nodes simultaneously when choosing the next move



# AND/OR Best-First-Search

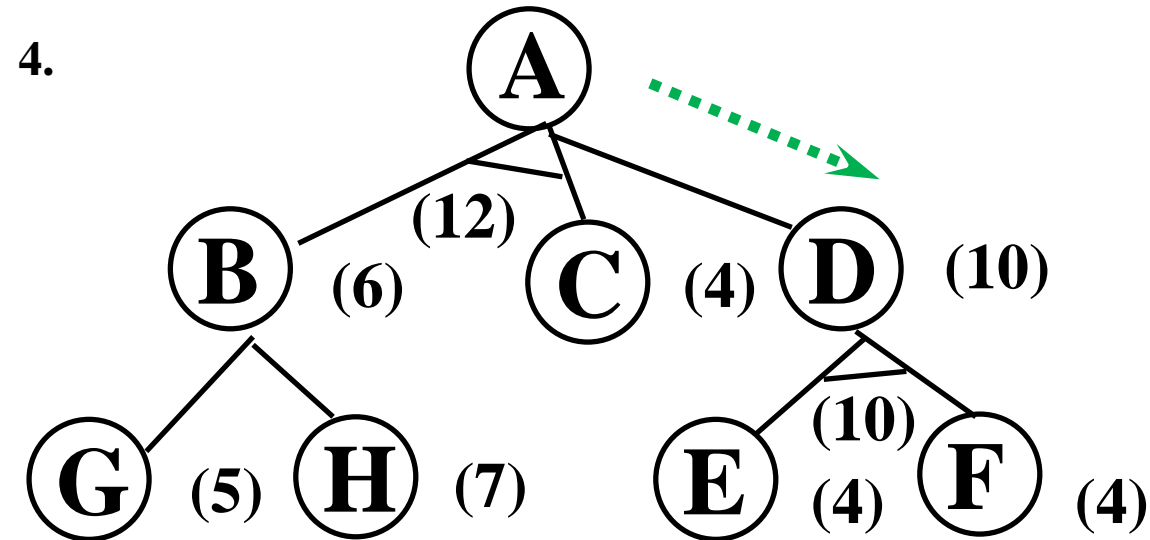
- Traverse the graph (from the initial node) following the best current path.
- Pick one of the unexpanded nodes on that path and expand it. Add its successors to the graph and compute  $f$  for each of them
- Change the expanded node's  $f$ -value to reflect its successors. Propagate the change up the graph.
- Reconsider the current best solution and repeat until a solution is found

# AND/OR Search example (Cont.)

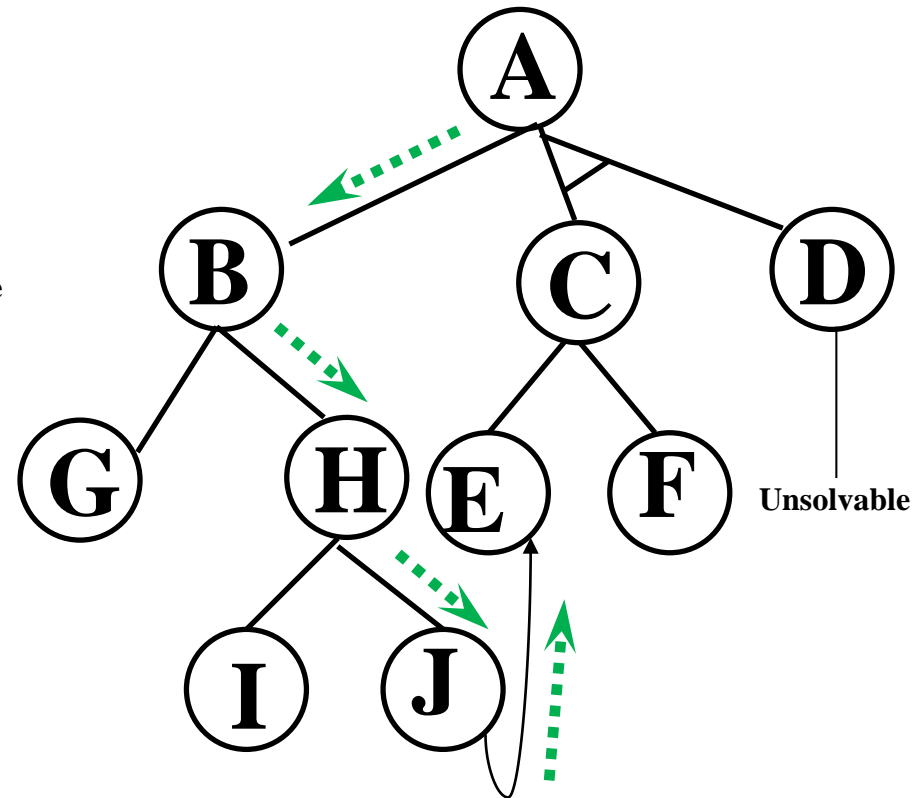
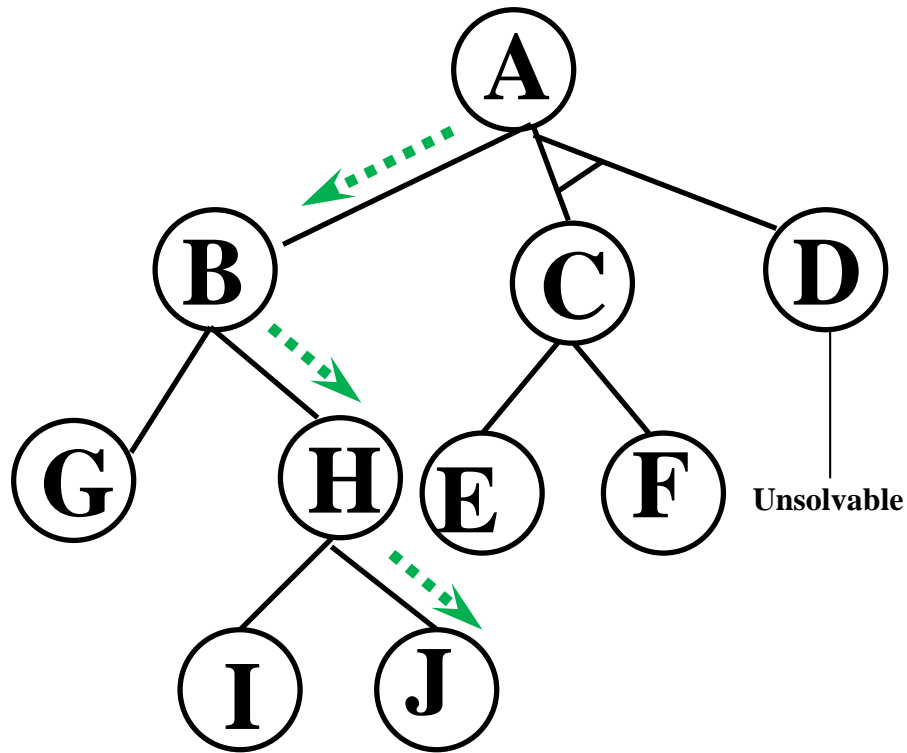




# AND/OR Search example

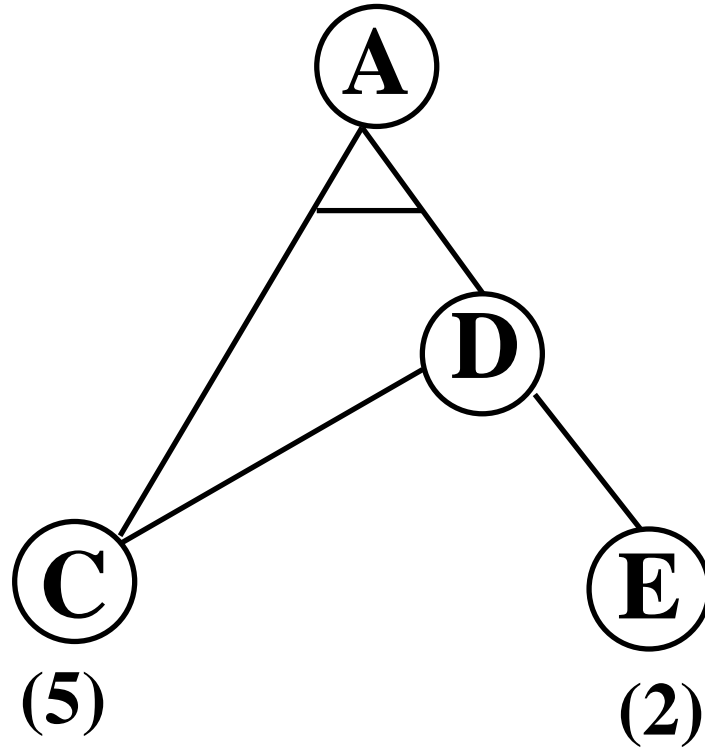


# A Longer Path May Be Better



# Interacting Sub Goals

---



# AO\* Algorithm

1. Let  $G$  be a graph with only starting node  $INIT$ .
2. Repeat the followings until  $INIT$  is labeled SOLVED or  $h(INIT) > FUTILITY$ 
  - a) Select an unexpanded node  $N$  from the most promising path from  $INIT$
  - b) Generate successors of  $N$ . If there are none, set  $h(N) = FUTILITY$ ; otherwise for each SUCCESSOR that is not an ancestor of  $N$ :
    - i. Add SUCCESSOR to  $G$ .
    - ii. If SUCCESSOR is a terminal node, label it SOLVED and set  $h(SUCCESSOR) = 0$ .
    - iii. If SUCCESSOR is not a terminal node, compute its  $h$

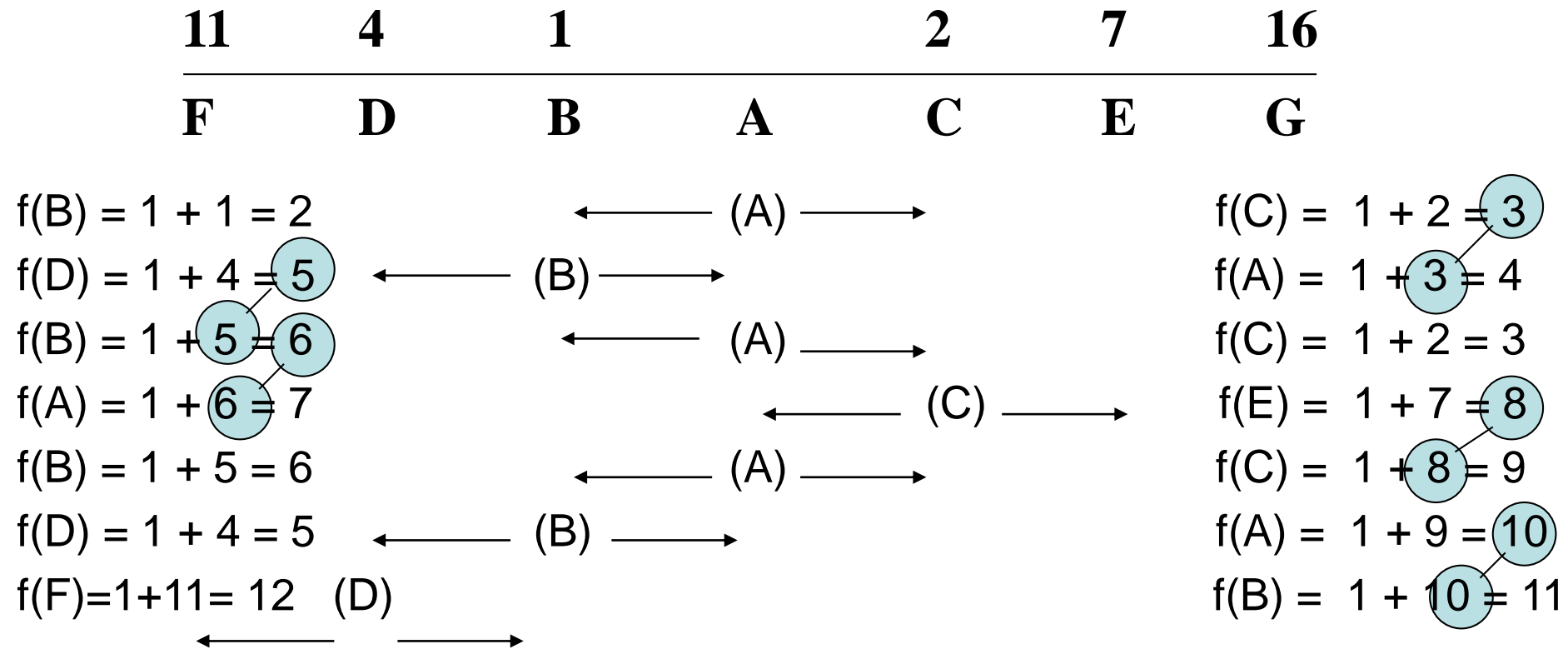
# AO\* Algorithm (Cont.)

- c) let  $S$  be set of SOLVED nodes or nodes whose  $h$  values have been changed and need to have values propagated back to their parents. Initialize  $S$  to  $N$ . Until  $S$  is empty repeat the followings:
- i. Remove a node from  $S$  and call it  $C$ .
  - ii. Compute the cost of each of the arcs emerging from  $C$ . Assign minimum cost of its successors as its  $h$ .
  - iii. Mark the best path out of  $C$  by marking the arc that had the minimum cost in step ii
  - iv. Mark  $C$  as SOLVED if all of the nodes connected to it through new labeled arc have been labeled SOLVED
  - v. If  $C$  has been labeled SOLVED or its cost was just changed, propagate its new cost back up through the graph. So add all of the ancestors of  $C$  to  $S$ .

# Real Time A\* (Online Search)

- *Backtracking* has a uniform (or actual) cost
- Considers the cost ( $> 0$ ) for switching from **one** branch to another (i.e., *backtracking*) in the search
- Uses a heuristic function similar to that of A\*
  - $f(n) = g(n) + h(n)$
- But,  $g$  is calculated from the current state, instead of the initial state

# Example: path finding in real life



# Example 2

Current State = S

$$f(A) = 3 + 5 = 8$$

$$f(B) = 2 + 4 = 6$$

Current State = B

$$f(S) = 2 + 8 = 10$$

$$f(A) = 4 + 5 = 9$$

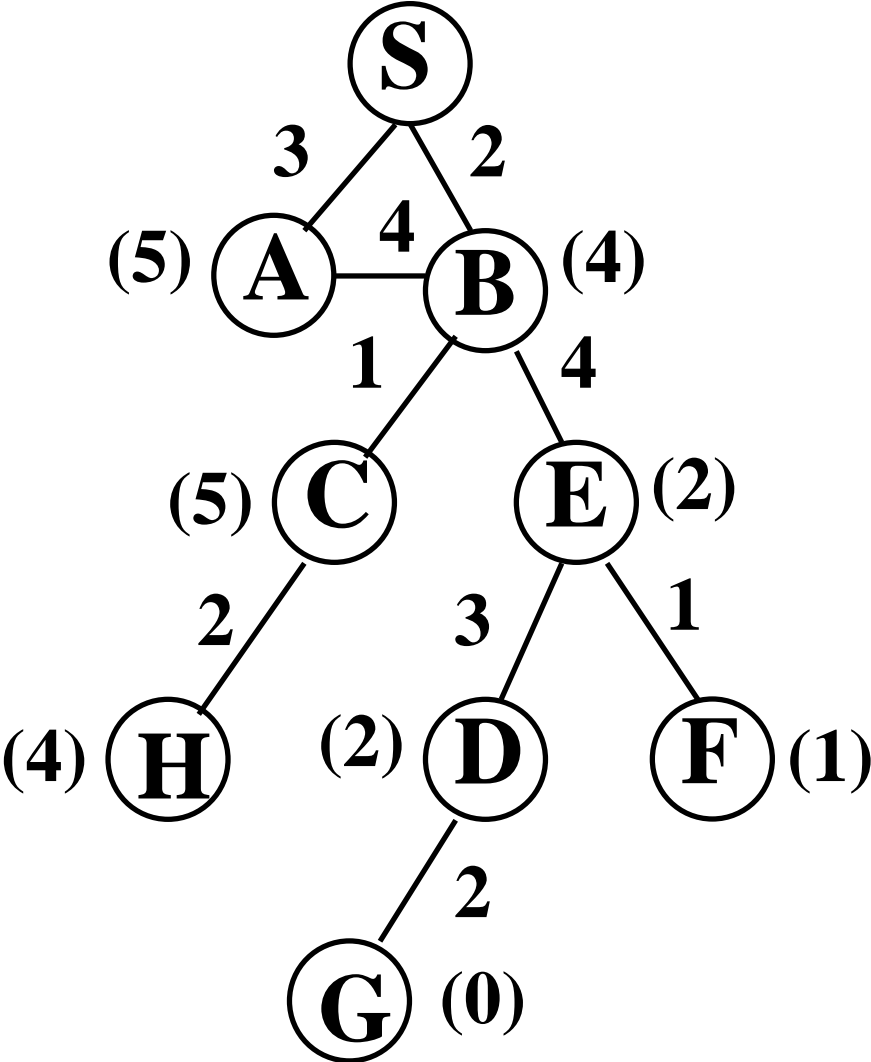
$$f(C) = 1 + 5 = 6$$

$$f(E) = 4 + 2 = 6$$

Current State = C

$$f(H) = 2 + 4 = 6$$

$$f(B) = 1 + 6 = 7$$





# Example 2

Current State = H

$$f(C) = 2 + 7 = 9$$

Current State = C

$$f(B) = 1 + 6 = 7$$

$$f(H) = \infty$$

Current State = B

$$f(S) = 2 + 8 = 10$$

$$f(A) = 4 + 5 = 9$$

$$f(E) = 4 + 2 = 6$$

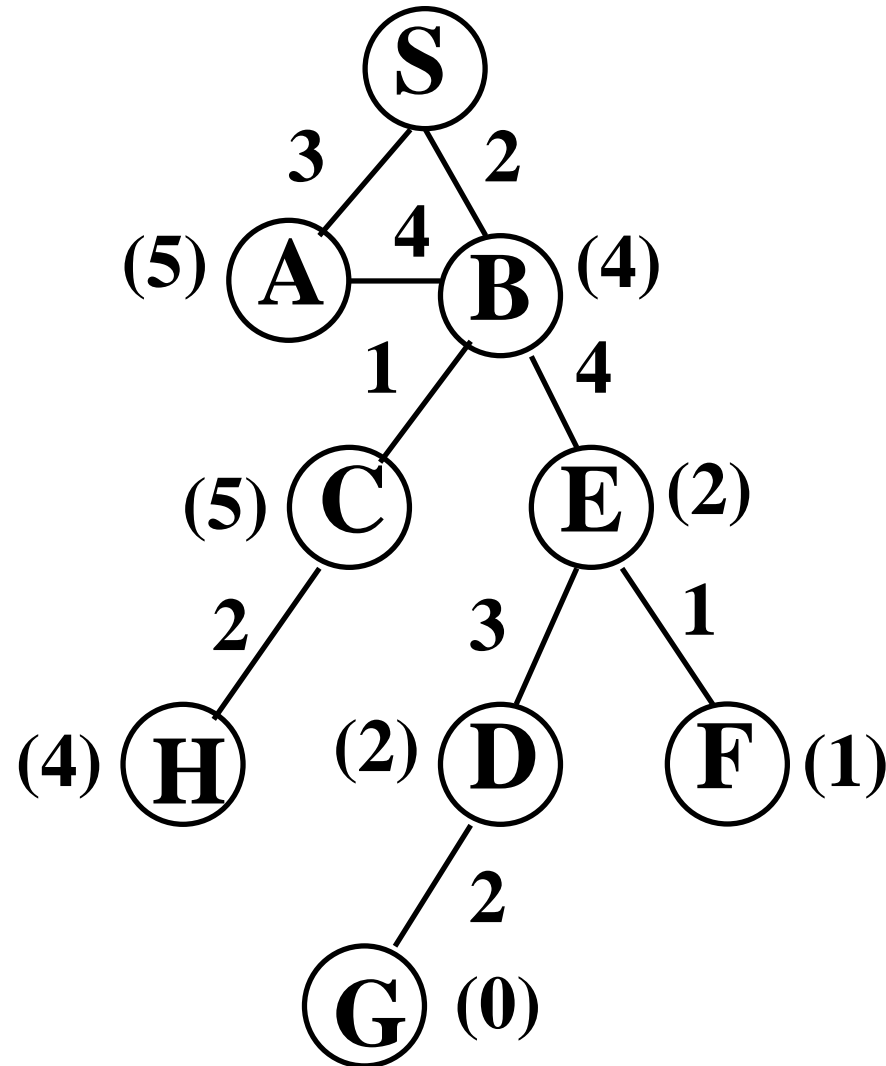
$$f(C) = \infty$$

Current State = E

$$f(B) = 4 + 9 = 13$$

$$f(D) = 3 + 2 = 5$$

$$f(F) = 1 + 1 = 2$$



# Example 2

Current State = F

$$f(E) = 2 + 5 = 7$$

Current State = E

$$f(D) = 3 + 2 = 5$$

$$f(B) = 4 + 9 = 13$$

$$f(F) = \infty$$

Current State = D

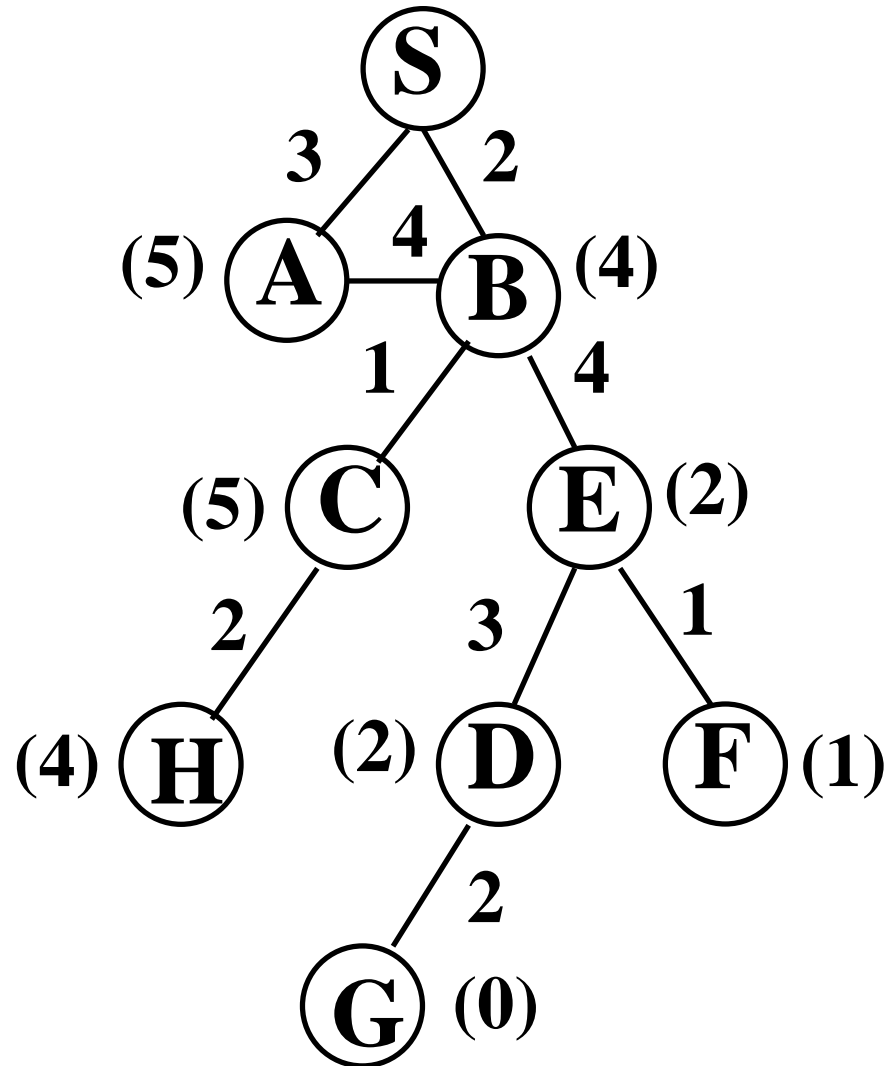
$$f(G) = 2 + 0 = 2$$

$$f(E) = 3 + 13 = 16$$

Visited Nodes =

S, B, C, H, C, B, E, F, E, D, G

Path = S, B, E, D, G



# Real Time A\* (Online Search)

- ◆ Maintains in a hash table a list of those states/nodes that have been visited by an *actual move* in the real world of the problem solver;
- ◆ At each cycle in the real-world, the current state is expanded and the heuristic function, possibly augmented by look-ahead search, is applied to each successor state which is **not in the hash table**;
- ◆ The  $f$  value of each neighboring state is computed by adding the  $h$  value plus the cost of the link to the current state;
- ◆ The neighbor with the minimum  $f$  value is chosen for the current state;
- ◆ The **second best  $f$  value is stored in the hash table** for the current state
  - Represents the estimated  $h$  cost of solving the problem by returning to this state
  - Second best avoids loops