# Look Into Details: The Benefits of Fine-Grain Streaming Buffer Analysis

Mohammad H. Foroozannejad Matin Hashemi Trevor L. Hodges Soheil Ghiasi

University of California, Davis, CA

## **Streaming Applications**

Widespread





Cell phones, mp3 players, video conference, real-time encryption, graphics, HDTV editing hyperspectral imaging, cellular base stations

#### Definition

- Infinite sequence of data items
- At any given time, operates on a small window of this sequence
- Moves forward in data space

//53° around the z axis const R[3][3]={ {0.6,-0.8, 0.0}, {0.8, 0.6, 0.0}, {0.0, 0.0, 1.0}} Rotation3D { for (i=0; i<3; i++) for (j=0; j<3; j++) B[i] += R[i][j] \* A[j] }

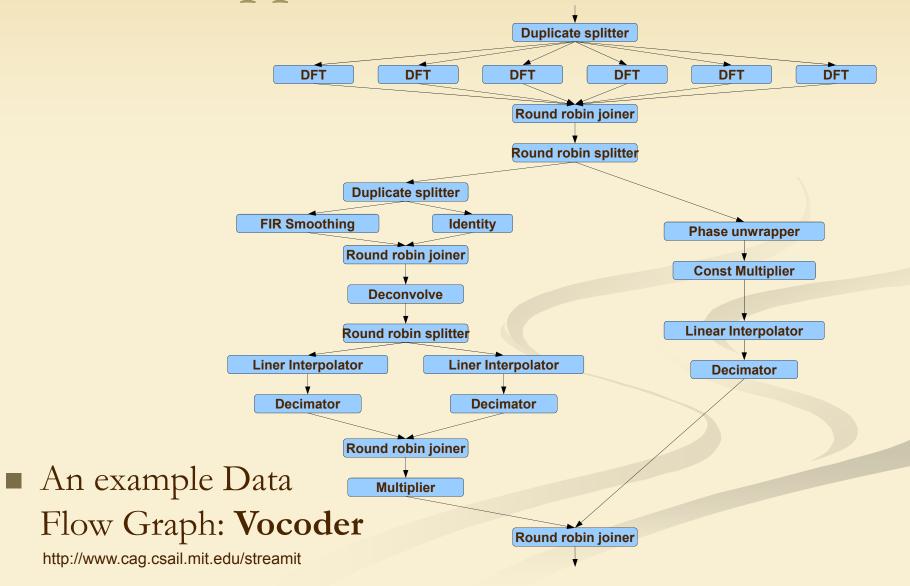


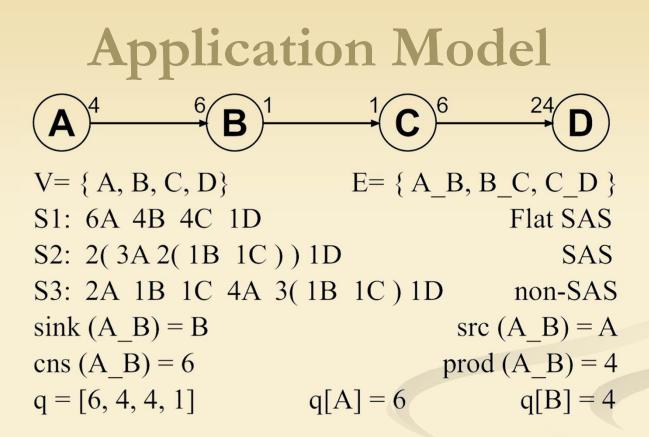
## **Application Model**

#### Data Flow Graph

- Vertices or Actors
  - functions, computations
- Edges
  - data dependency, communication between actors
- Execution Model
  - any actor can perform its computation whenever all necessary input data are available on incoming edges.

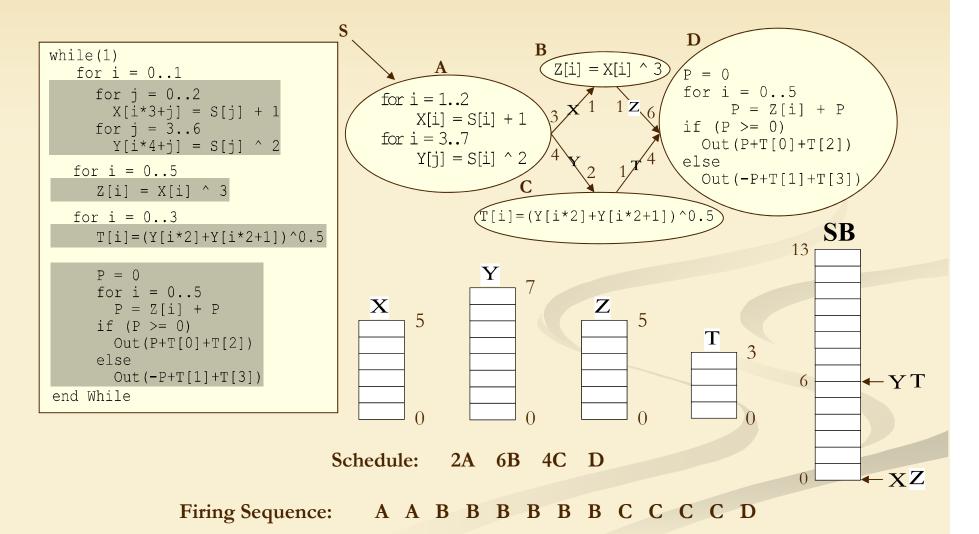
### **Application Model**

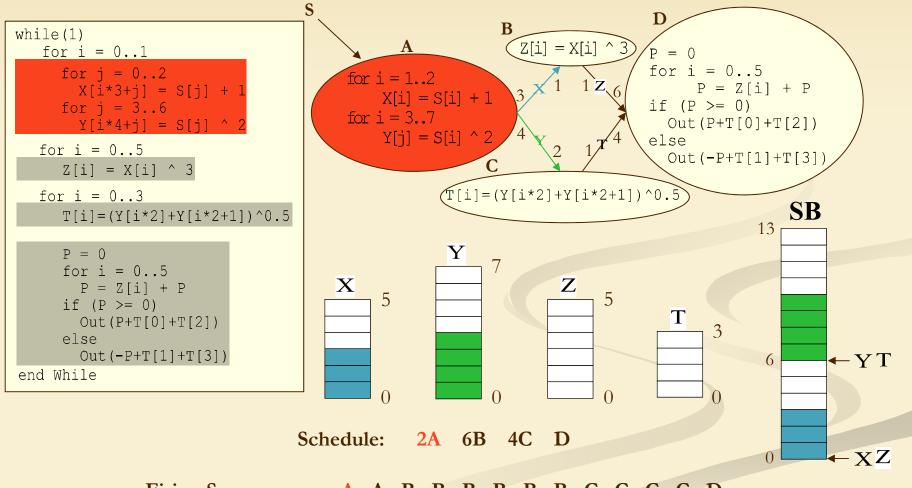


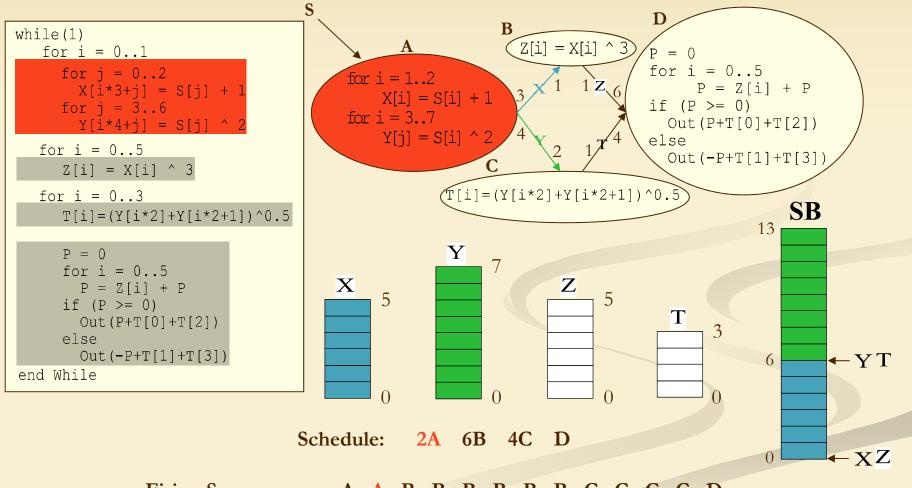


SDF (Synchronous Data Flow Graph) is one special case

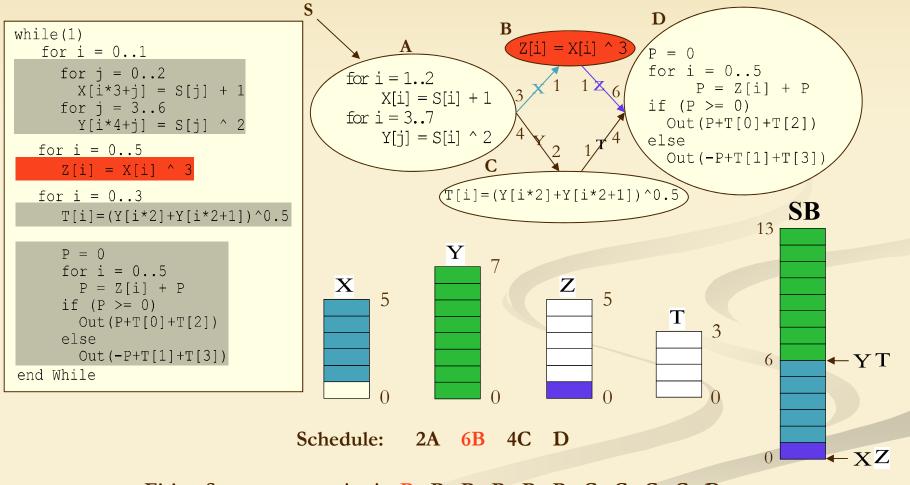
- Fixed input and output rates on the edges
- statically schedulable



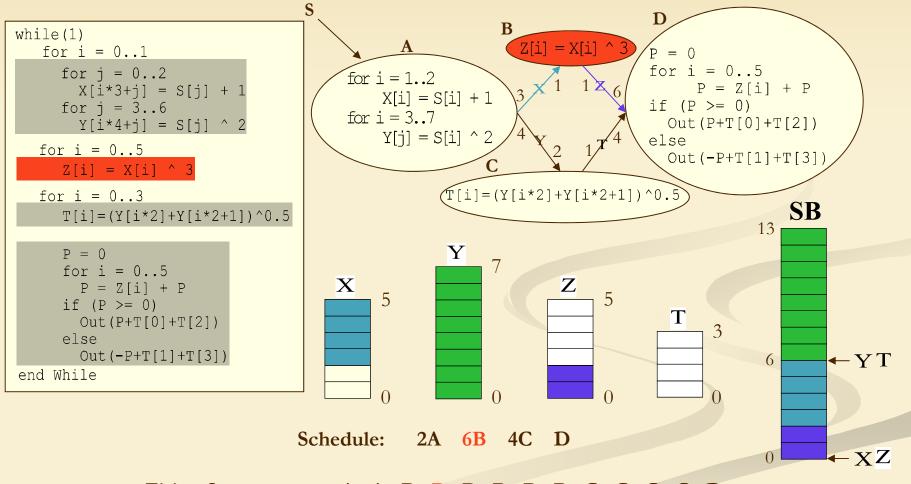


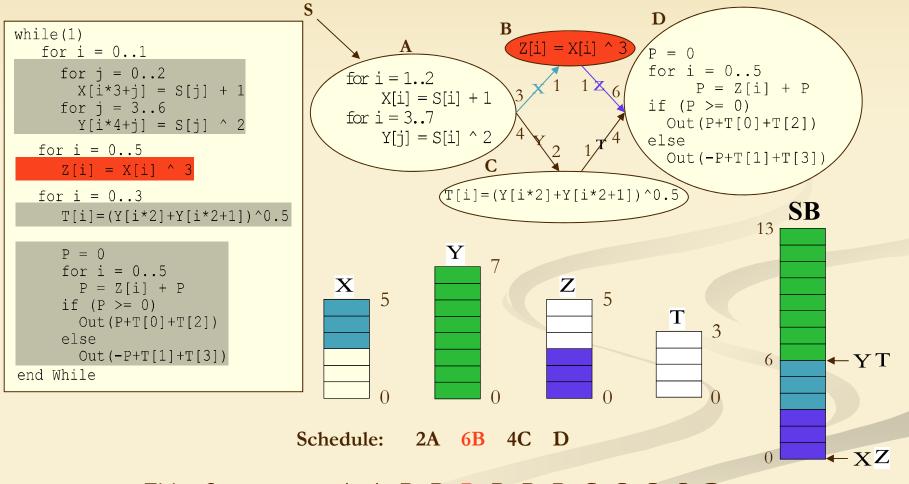


Firing Sequence: A A B B B B B B C C C C D

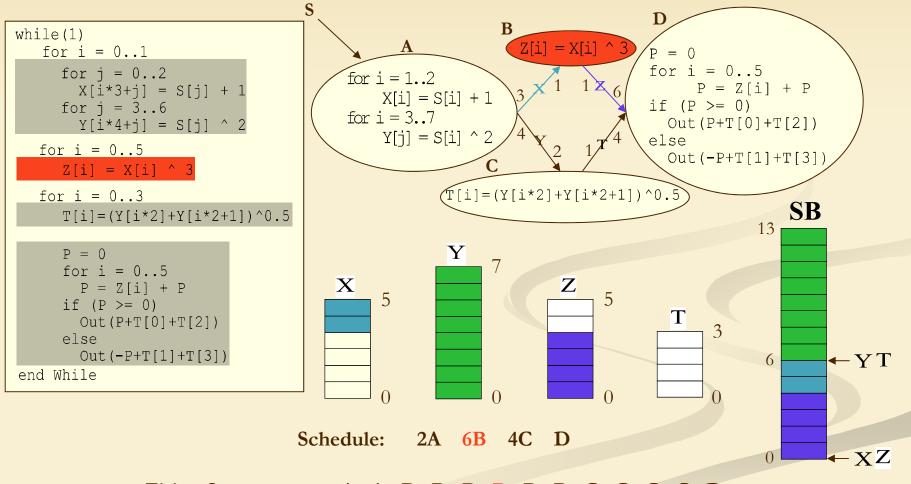


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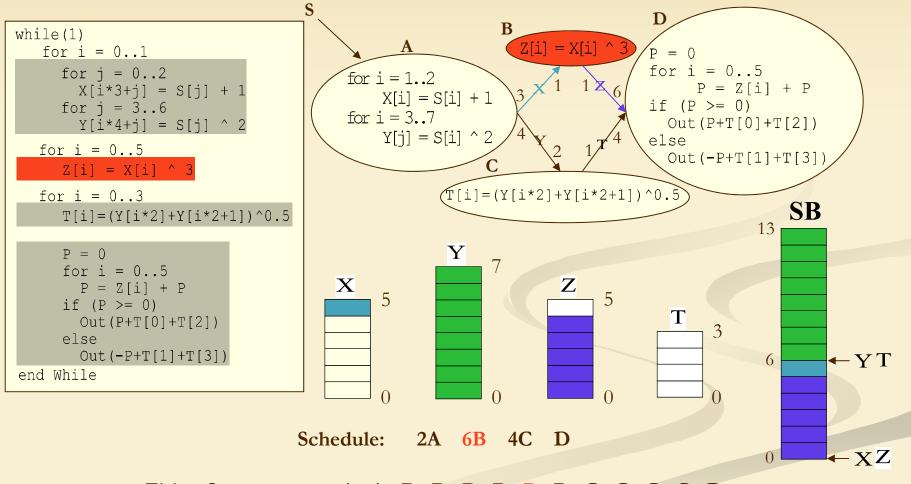




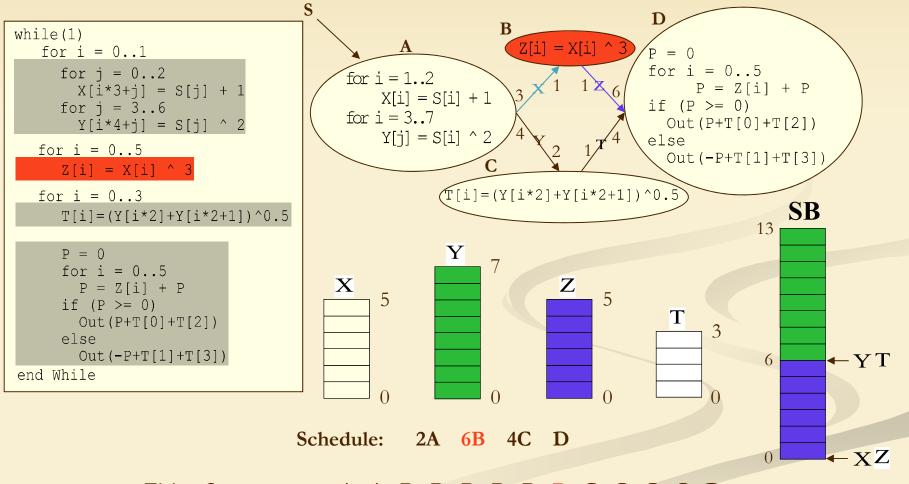
Firing Sequence: A A B B B B B B C C C C D



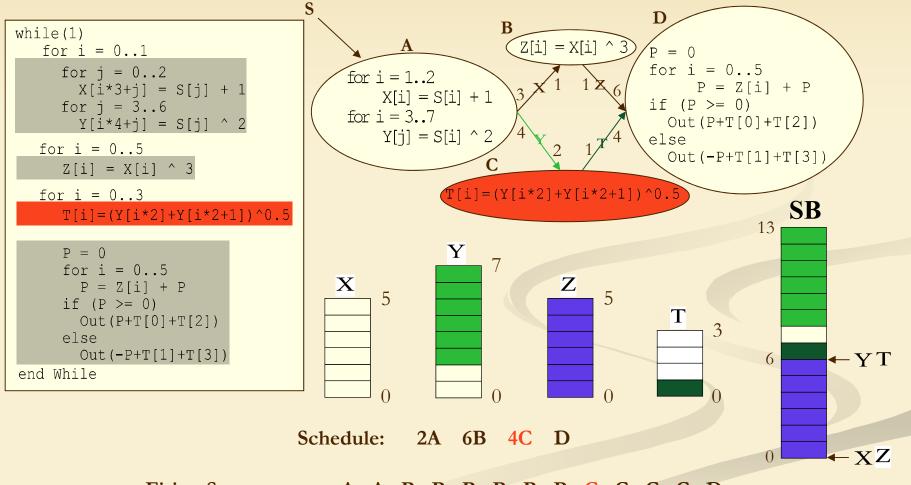
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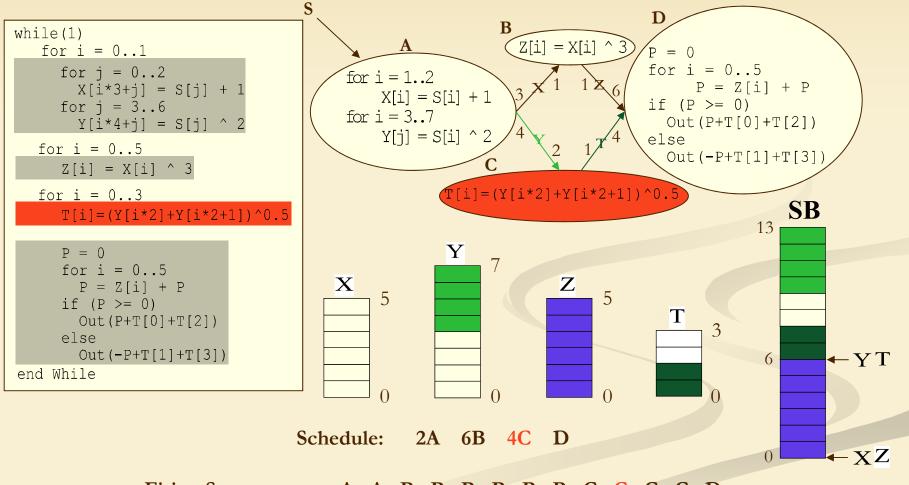
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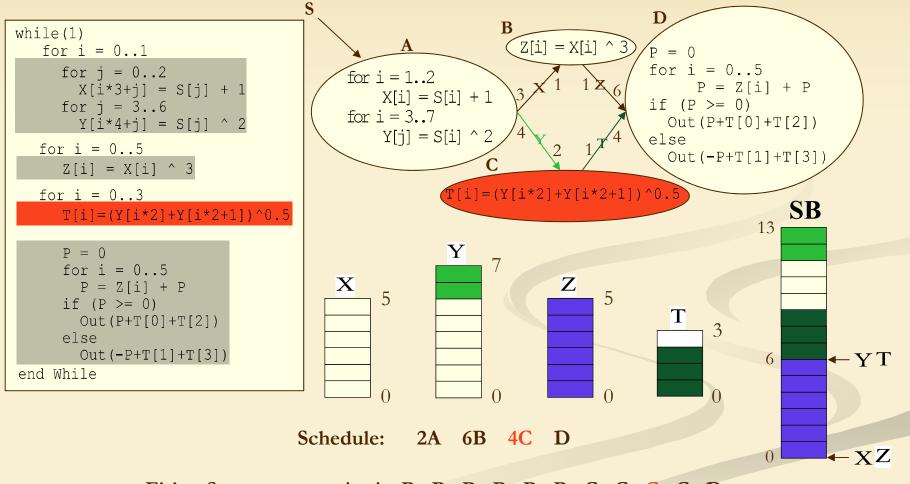
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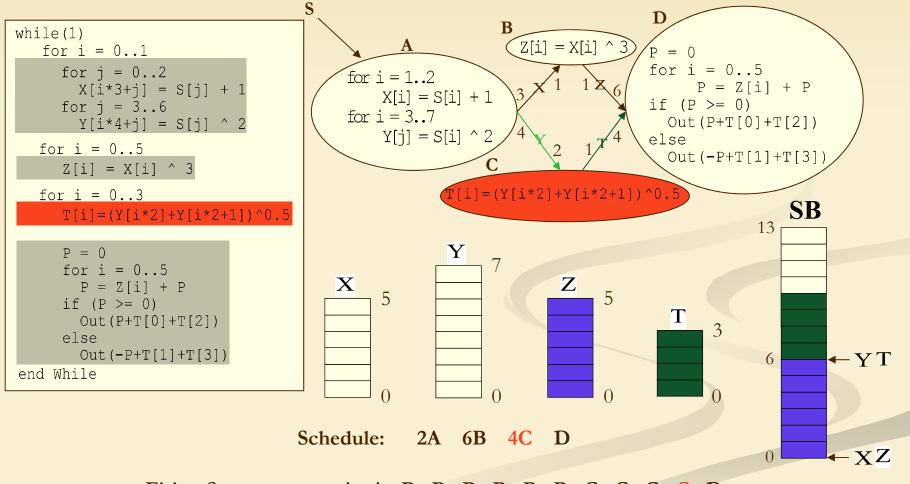
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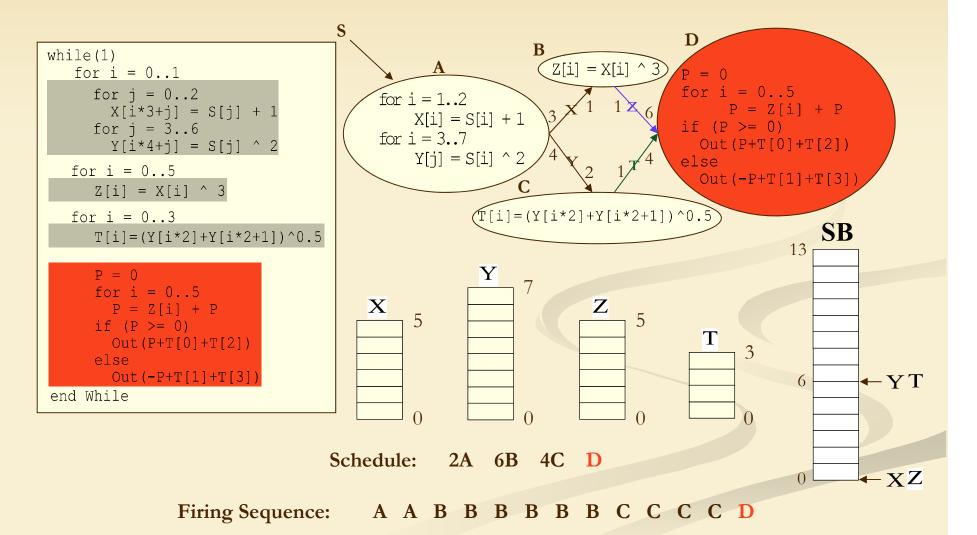
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Firing Sequence: A A B B B B B B C C C D



Firing Sequence: A A B B B B B B C C C C D



## **Shared Buffer Implementation**

#### Idea:

- Most of the time channel buffers are completely or partially empty.
- Rules:
  - 1. No over-writing or reading another buffer's data.
  - 2. Statically allocated
  - 3. No re-allocation

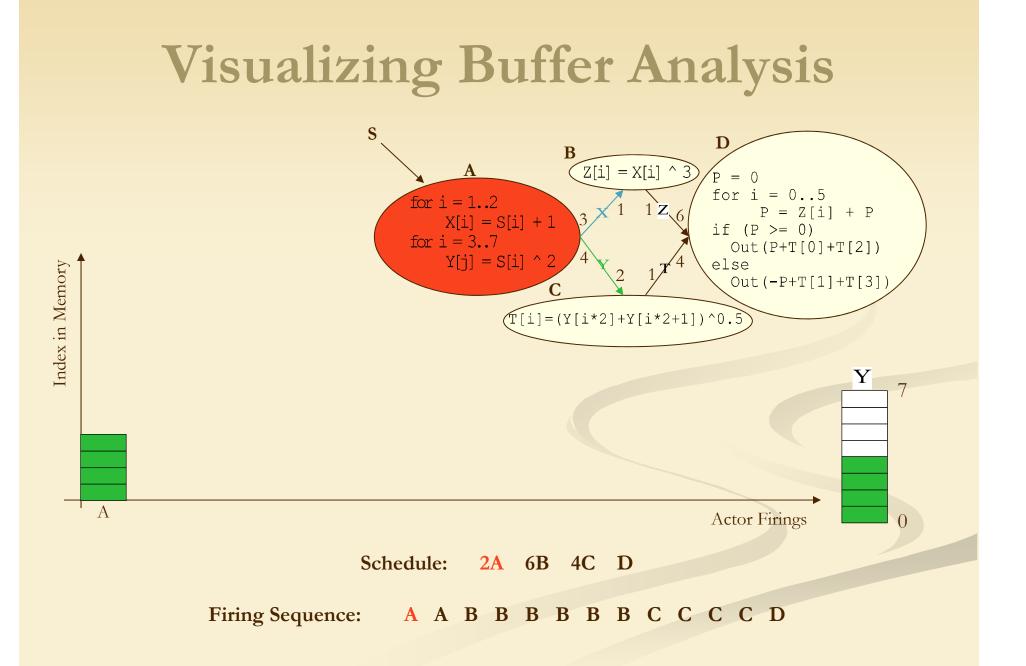
## Visualizing Buffer Analysis

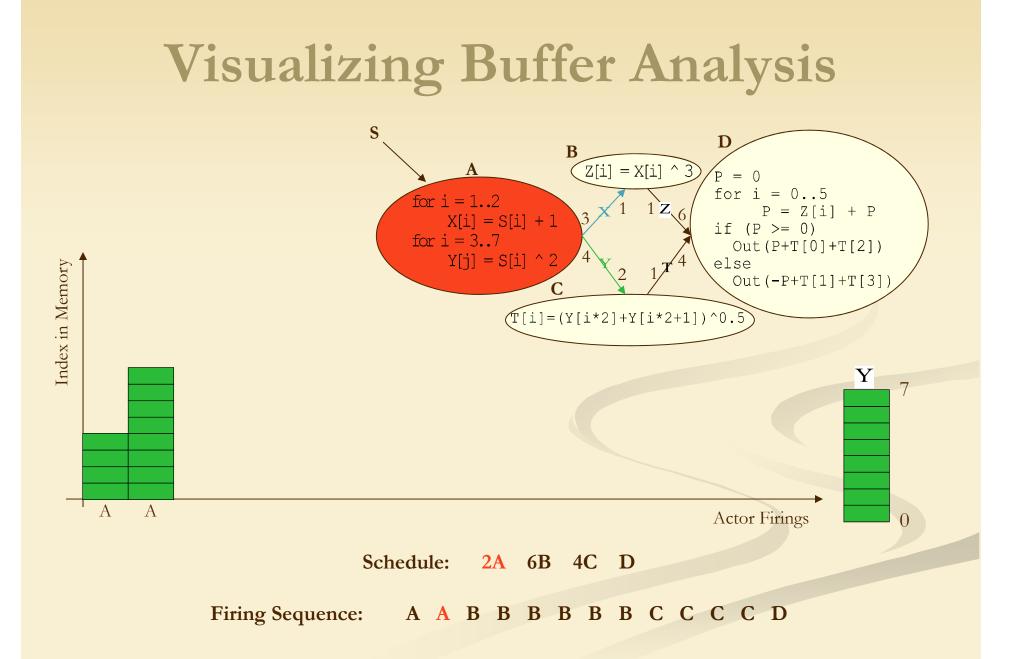
#### Tow dimensional plane

- X-axis: Actor firings in the schedule (time)
- Y-axis: Buffer location in the memory (space)
- Filled Area: The range between Head and Tail indices

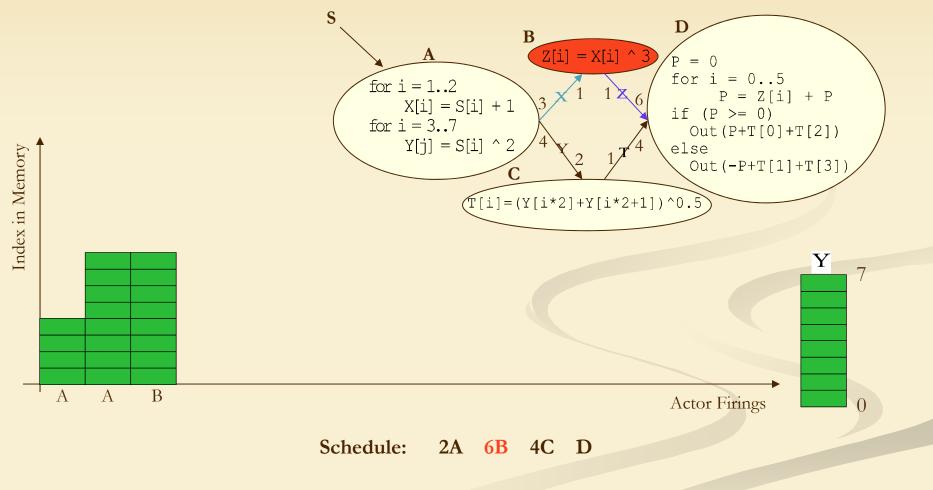
#### Advantage:

- Memory allocation problem can be viewed as a geometric layout instance
- A solution is valid when the laid out buffers do not conflict in the time-memory plane.

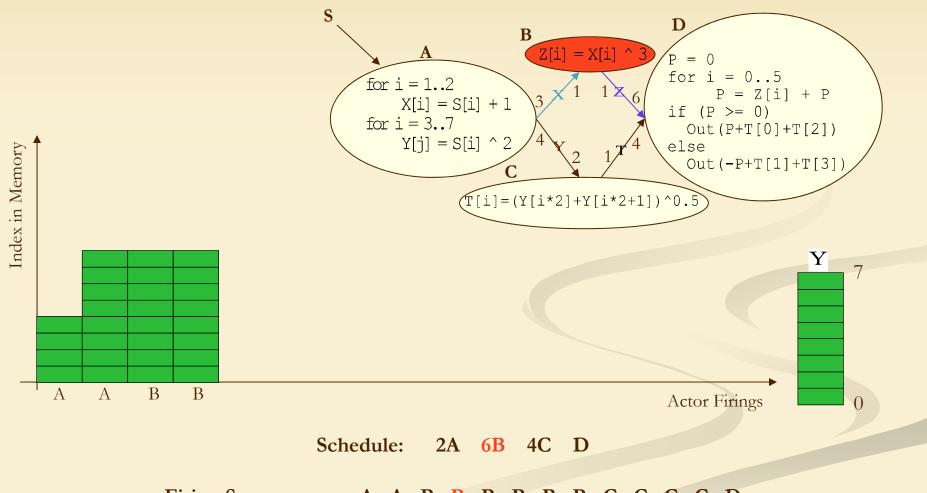






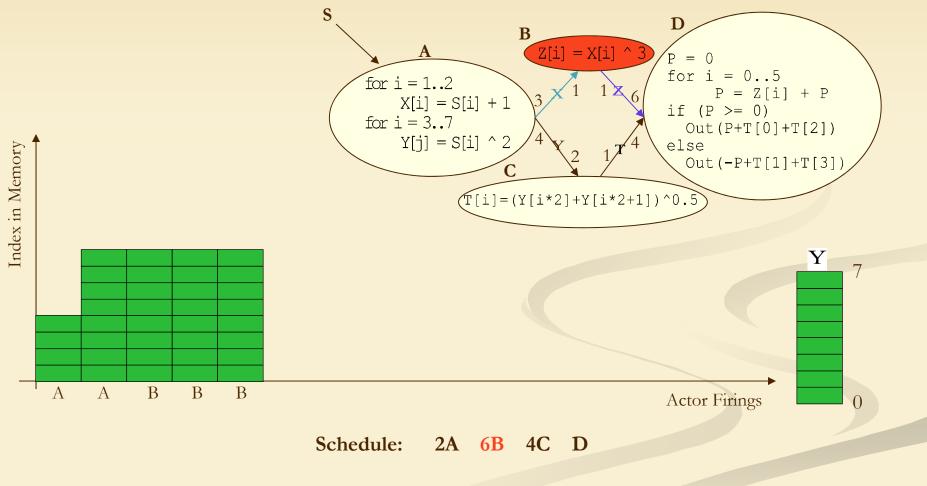




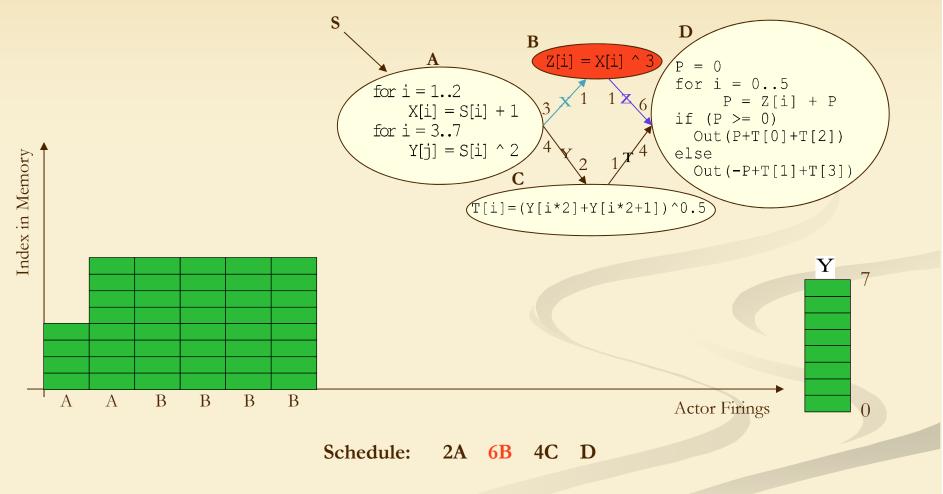


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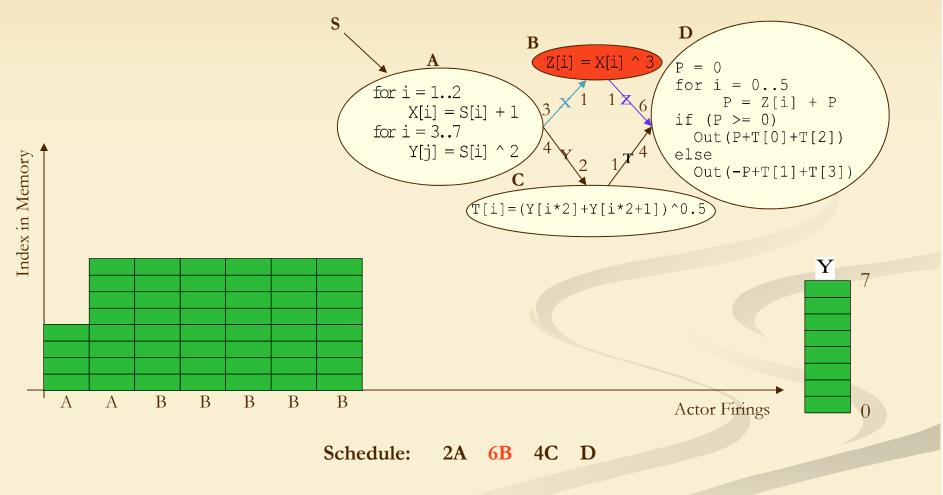






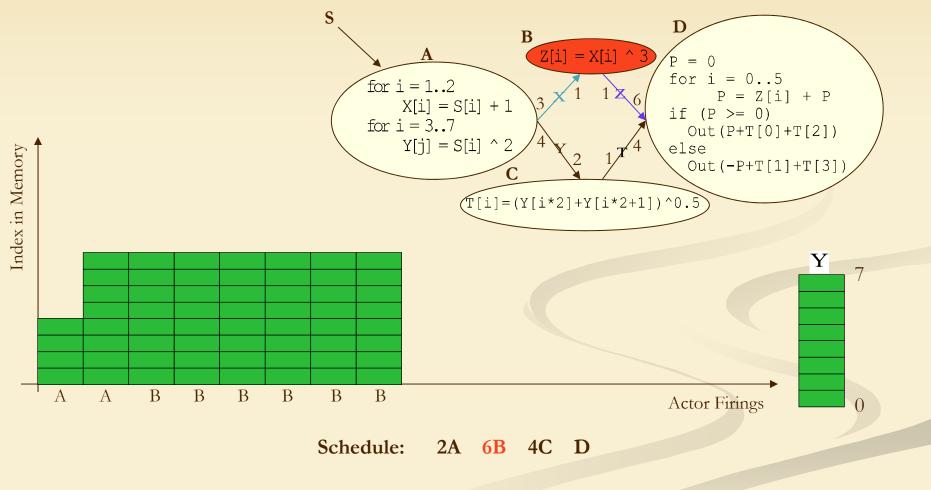
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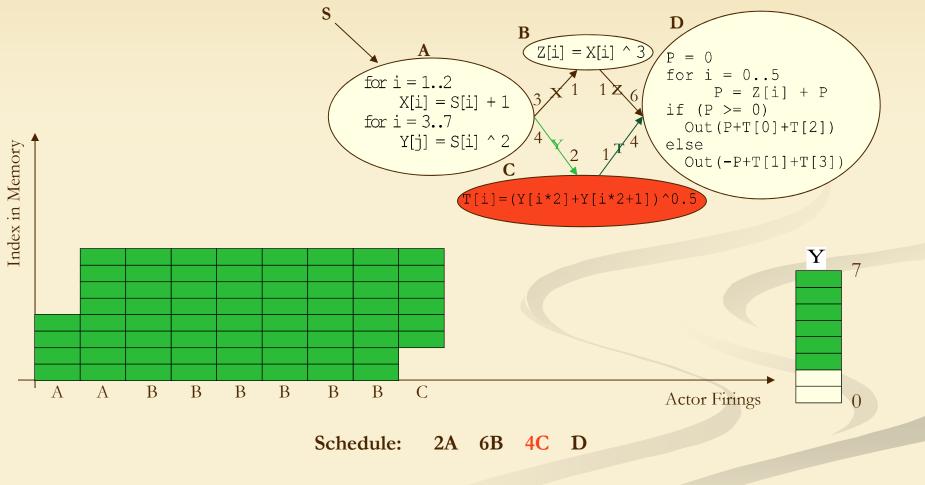


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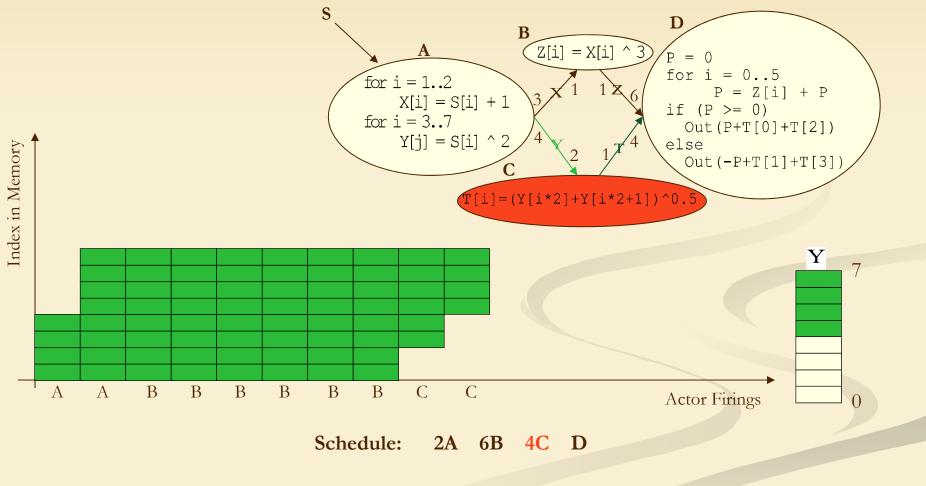




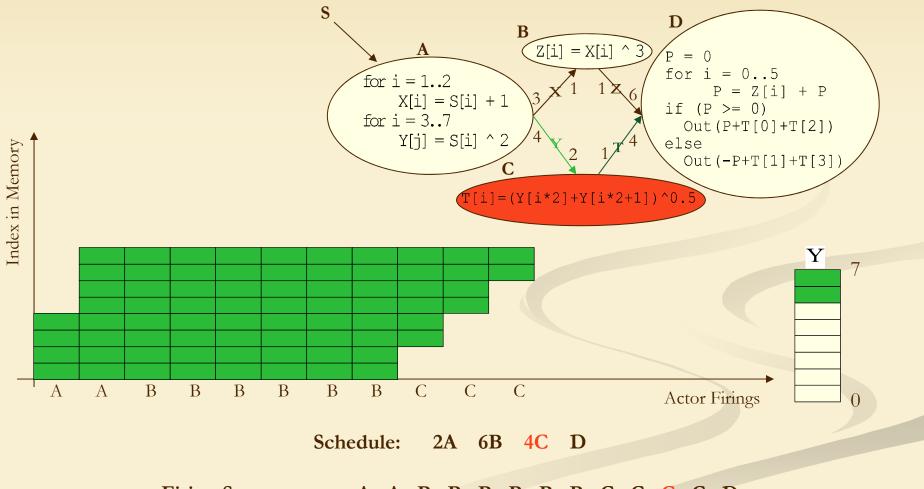


, sequence.

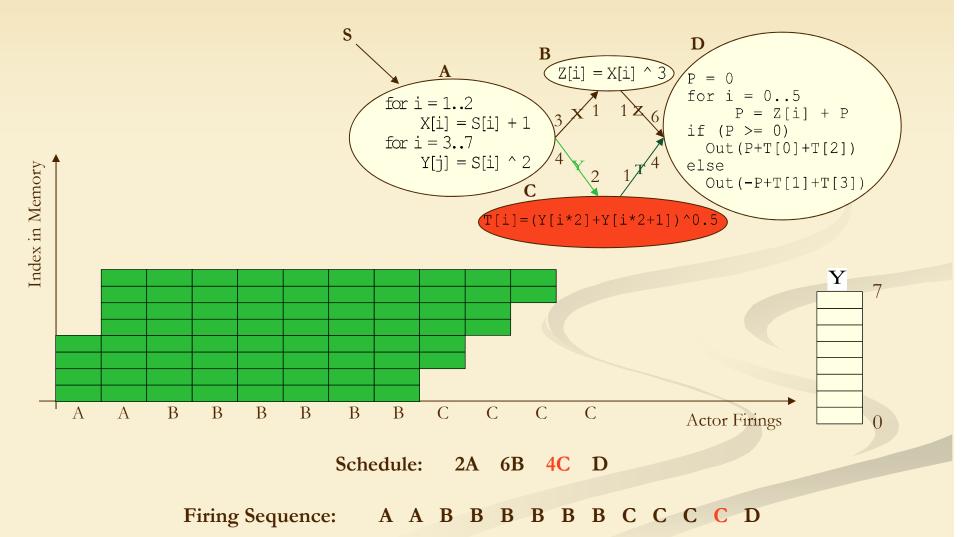




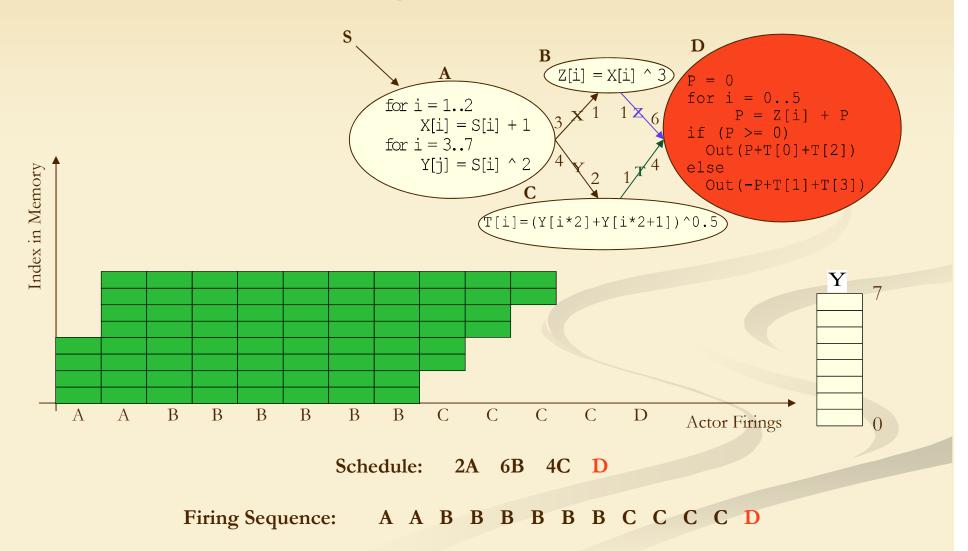






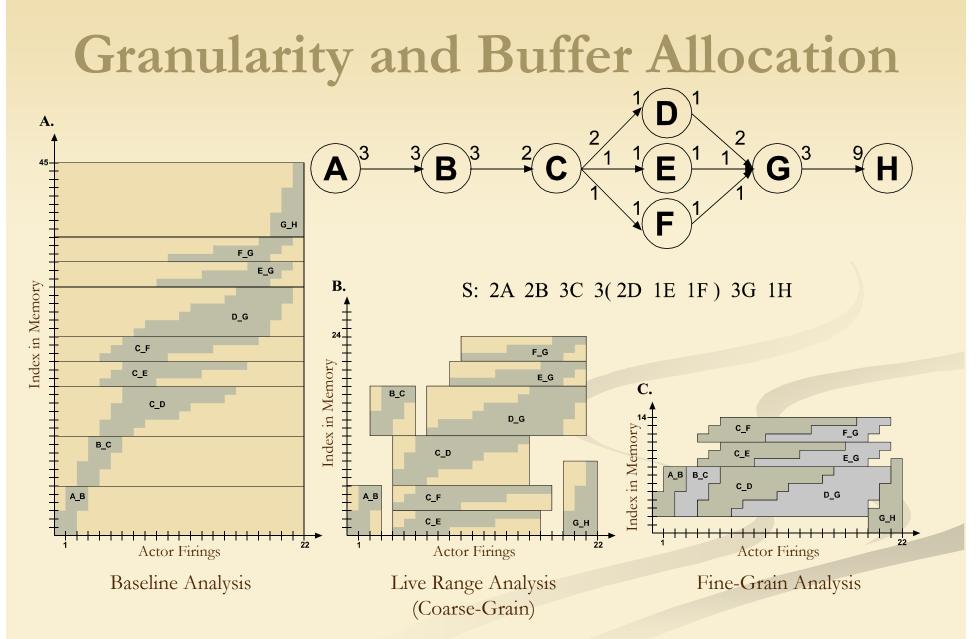


## Visualizing Buffer Analysis



### **Granularity and Buffer Allocation**

- The granularity in buffer analysis compromises accuracy in temporal behavior of buffers with analysis complexity:
  - Baseline
  - Coarse-grain
  - Fine-grain



P. K. Murthy and S. S. Bhattacharyya. Shared buffer implementations of signal processing systems using lifetime analysis techniques.

#### **Fine-Grain Buffer Allocation**

#### Mathematic Formulation:

- Use of existing tools
- Choose the best data structure

 $\forall e \in E : B_e = (H_e, L_e)$ 

 $B_e : \text{The Buffer on edge } e \text{ which we call it buffer } e \text{ in short}$  $H_e[t] : \text{Head index at time } 0 \le t \le T \text{ for the buffer on } e$  $L_e[t] : \text{Tail index at time } 0 \le t \le T \text{ for the buffer on } e$  $T = \sum_{v \in q_G} q[v]$ 

$$O = \{ (o_{e_1}, o_{e_2}, o_{e_3}, \dots, o_{e_N}) \mid e_1 : e_N \in E , N = |E| \}$$

#### **Fine-Grain Buffer Allocation**

#### • LEMMA:

In SA schedules the head index at the time t is always greater than equal the tail index at the same time:  $\forall t \leq T : H_e[t] \geq L_e[t]$ 

Constraints:

 $\forall e, b \in E \quad \forall 0 \le t \le T :$  $H_e[t] + o_e \le L_b[t] + o_b \quad OR \quad H_b[t] + o_b \le L_e[t] + o_e$ 

• Objective: Minimize <u>Shared Buffer Size</u>:

 $SBS = \max_{\forall e \in E} \{ o_e + H_e^{max} \mid H_e^{max} = \max_{0 \le t \le T} (H_e[t]) \}$ 

#### **ILP Formulation**

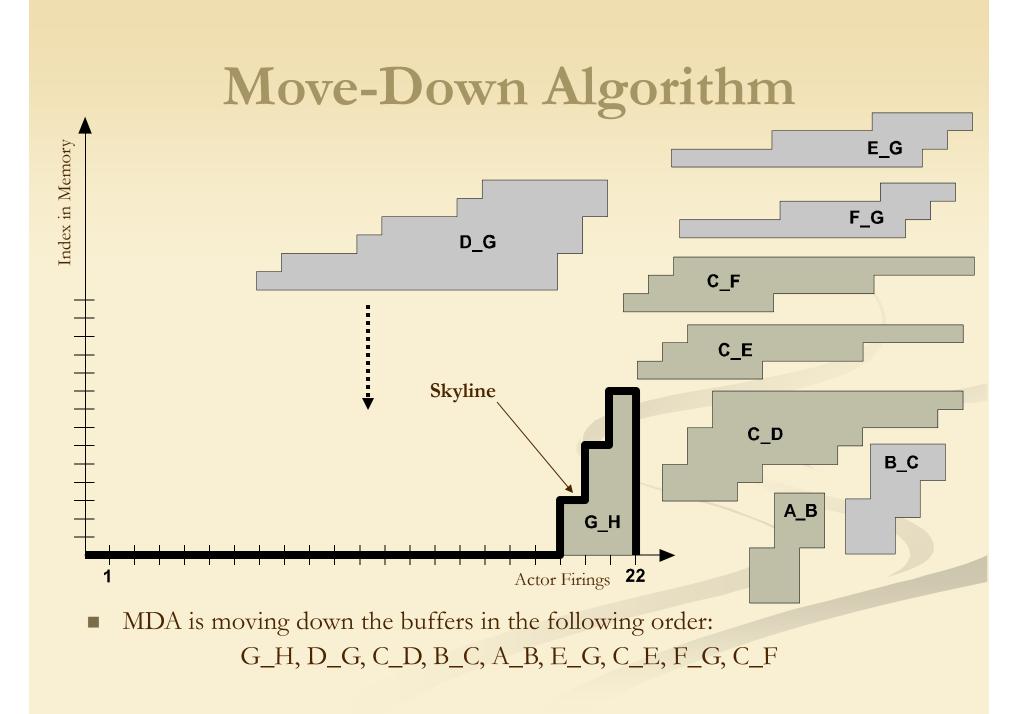
- The complexity of buffer sharing instance, and ILP runtime grows exponentially.
- Linear constraints cannot be easily used to articulate the "OR" logic:
  - Binary variables For each buffer and each location in the shared memory space
  - Constraints have to be generated for all time steps.

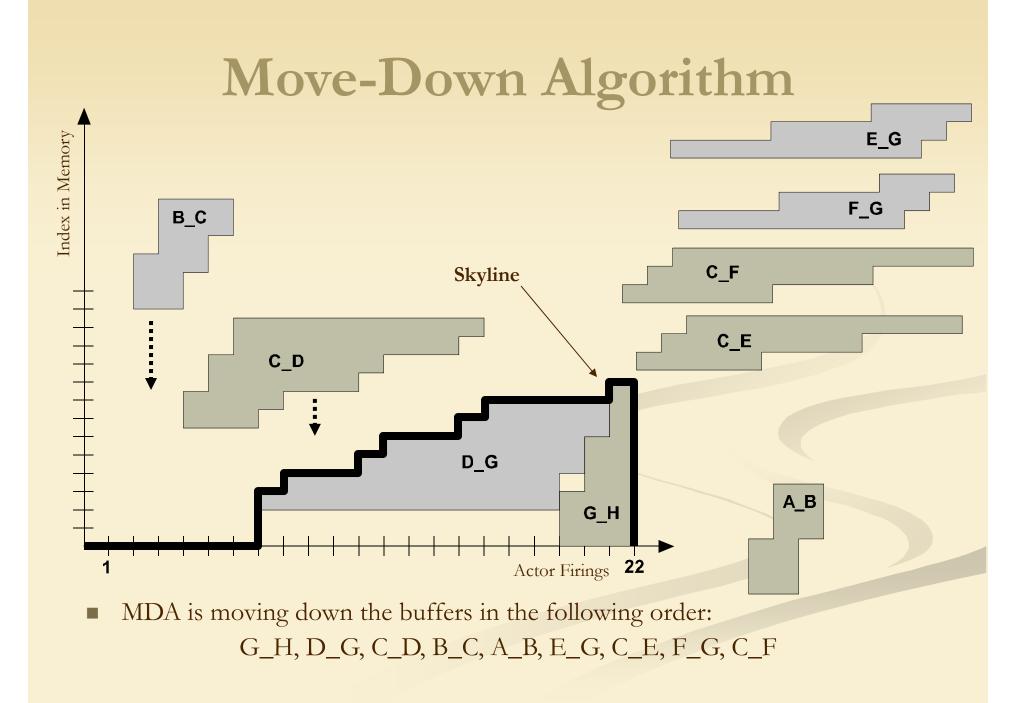
# Strip Packing Problem and Buffer-Sharing

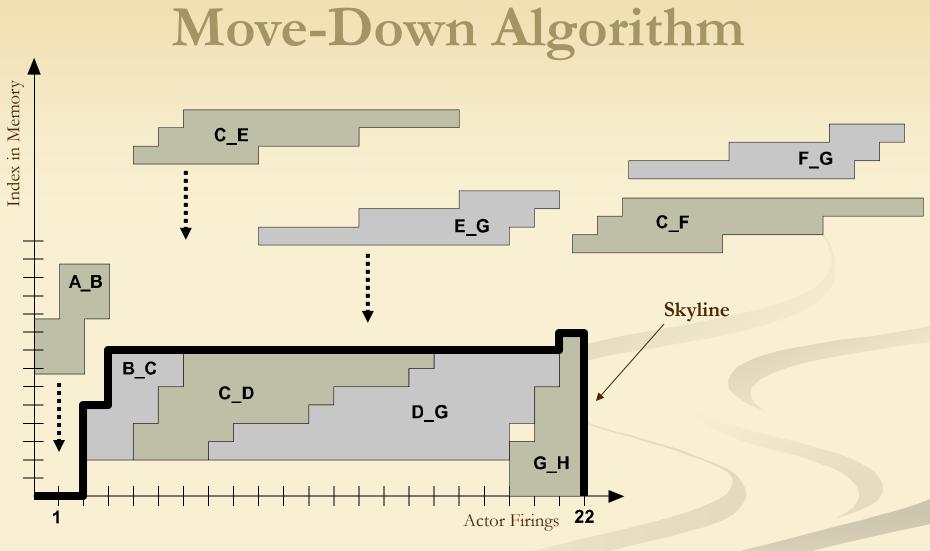
- In several industries there is a need for packing a set of 2-dimensional objects on a larger rectangular unit of material by minimizing the waste.
  - Two-Dimensional Bin Packing Problem (2BP):
    - wood or glass industries, warehousing contexts, newspapers paging
  - Two-Dimensional Strip Packing Problem (2SP):
    - paper or cloth industries

# Strip Packing Problem and Buffer-Sharing

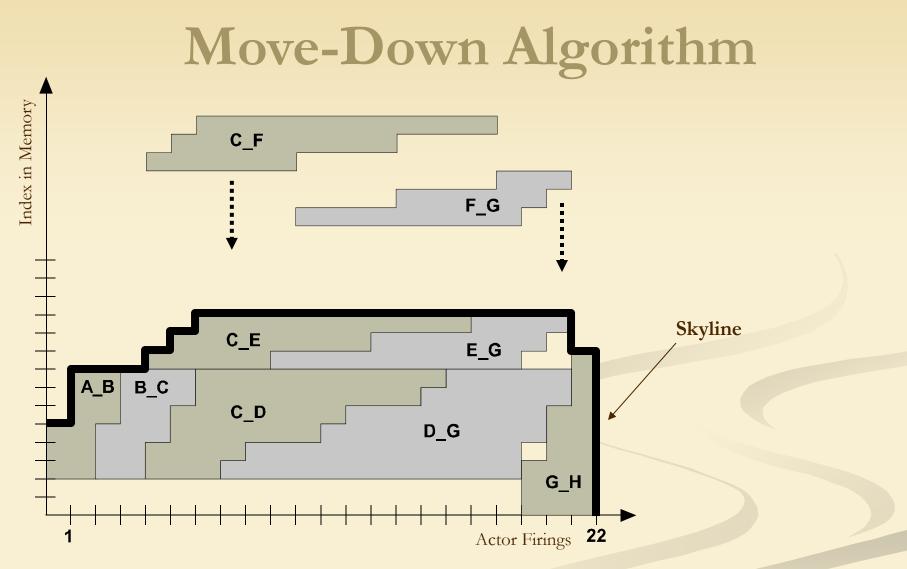
- The relationship between Packing Problems and Buffer Sharing Problem:
  - Objects: Buffer Size in Time which form complex polygons
  - Roll of Material: Shared Buffer Memory
  - Objective: To allocate an index to each buffer in the shared memory with no conflict using minimum space
  - Difference: We cannot move the objects (polygons) in time.
     We are only allowed to move them vertically. We also have no rotation.



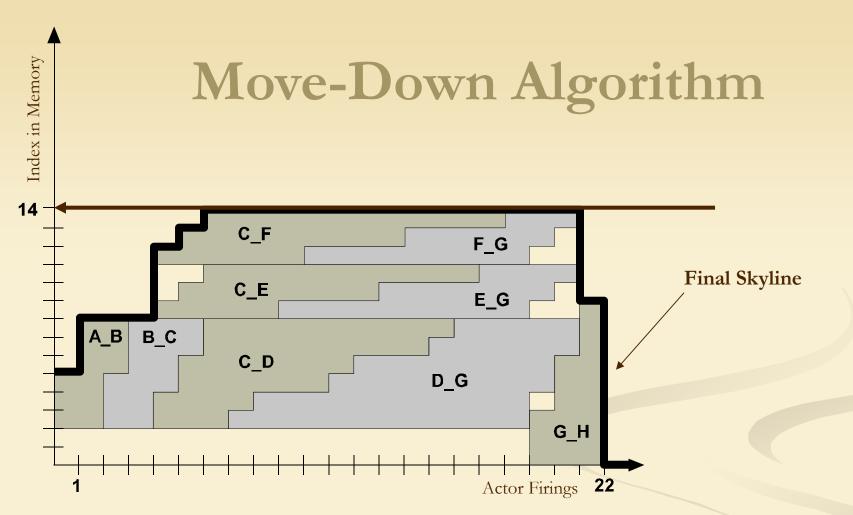




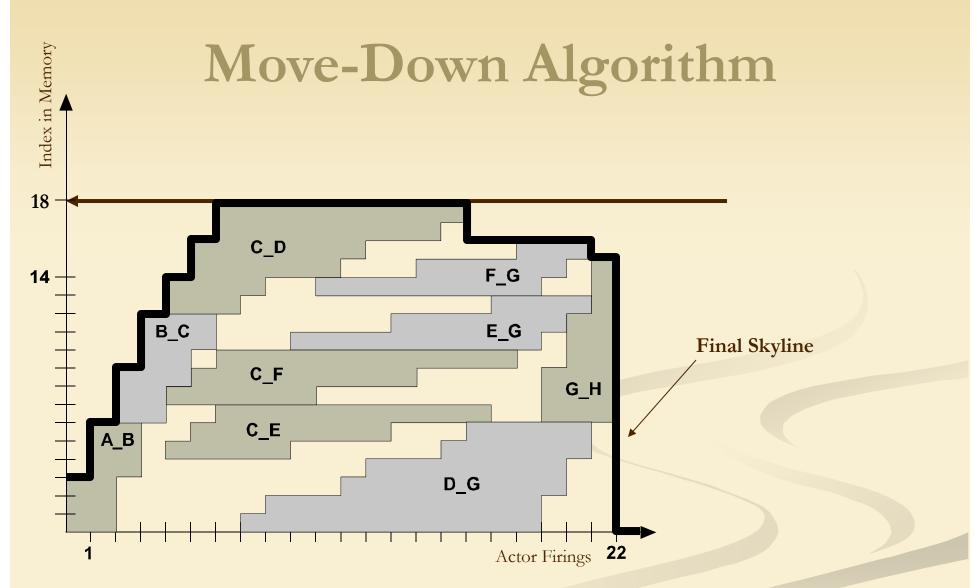
MDA is moving down the buffers in the following order:
 G\_H, D\_G, C\_D, B\_C, A\_B, E\_G, C\_E, F\_G, C\_F



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 G\_H, D\_G, C\_D, B\_C, A\_B, E\_G, C\_E, F\_G, C\_F



- The final placement of the buffers corresponding to the following order: G\_H, D\_G, C\_D, B\_C, A\_B, E\_G, C\_E, F\_G, C\_F
- The height of the final skyline indicates the shared memory size.



Another sequence which leads to the size 18 (14 is the optimal):
 A\_B, D\_G, G\_H, C\_E, C\_F, B\_C, E\_G, F\_G, C\_D

# Evolutionary Optimization using MDA

#### Genetic Algorithms in General:

- Chromosome: Provides an abstract representation of solutions in the search space,
- Inheritance: Models the basic operations through which, chromosomes are perturbed to improve the solution quality
  - Crossover
  - Mutation
- Fitness Function: Quantizes the quality of candidate solutions, and determines survival of selected candidates.

# Evolutionary Optimization using MDA

Initialization: Randomly select a set of permutations Sample set = { $\pi_1, \pi_2, \pi_3, \dots, \pi_N$  }

Fitness function: 
$$f(\pi) = \frac{1}{height(\pi)}$$

• Selection:  

$$p(\pi_i) = \frac{f(\pi_i)}{\sum_{j=1}^N f(\pi_j)}$$

## Evolutionary Optimization using MDA

#### Crossover:

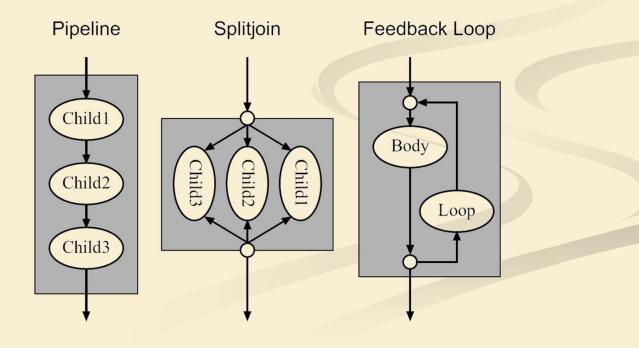
• Example: 
$$p = 2$$
  $q = 4$   
 $\pi_{parent1} = (B_{e1}, \underline{B}_{e2}, \underline{B}_{e3}, \underline{B}_{e4}, B_{e5}, B_{e6})$   
 $\pi_{parent2} = (\underline{B}_{e6}, \underline{B}_{e5}, B_{e4}, B_{e3}, B_{e2}, \underline{B}_{e1})$   
 $\pi_{child} = (B_{e2}, B_{e3}, B_{e4}, B_{e6}, B_{e5}, B_{e1})$ 

Mutation:

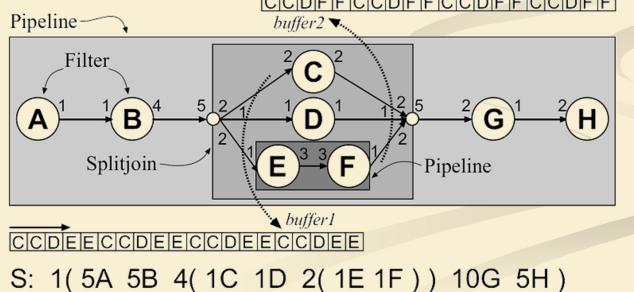
• Example: 
$$p_{mutation} = 0.4$$
 : the probability of being mutated  
 $i = 2$   $j = 4$   
 $\pi_{child}$  Before  $= (B_{e2}, \mathbf{B}_{e3}, B_{e4}, \mathbf{B}_{e6}, B_{e5}, B_{e1})$   
 $\pi_{child}$  After  $= (B_{e2}, B_{e6}, B_{e4}, B_{e3}, B_{e5}, B_{e1})$ 

Iteratively, new children are generated and compared to the existing members until the termination point where we can return the best solution found.

- We have integrated our algorithm into the MIT StreamIt compiler
- Three composite stream objects in StreamIt
- Filters specify data processing



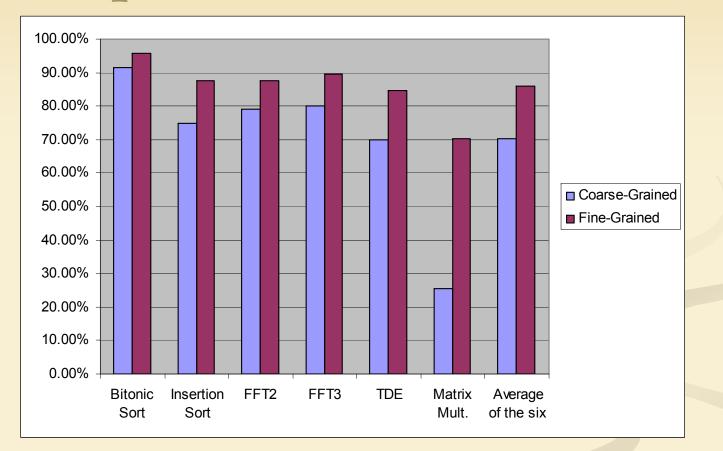
- The StreamIt scheduler is designed based on the hierarchical nature of the language.
- In Split-joins, one large buffer is used to implement multiple channels that either split to or join from several actors.



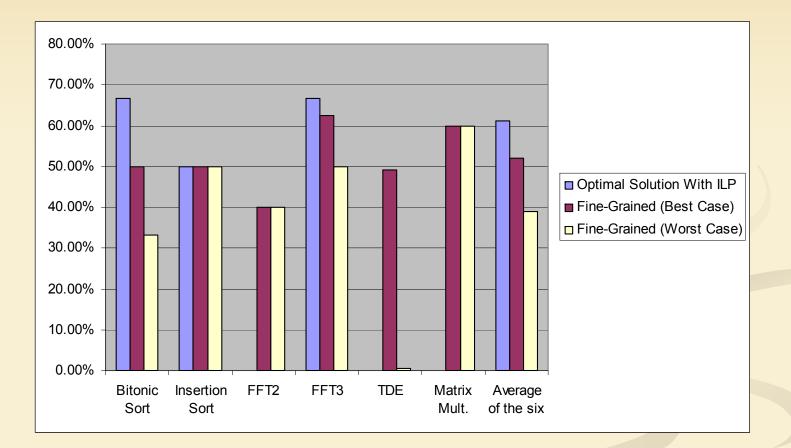
#### Benchmark Applications:

- Two sorting algorithms: Bitonic Sort, Insertion Sort
- Two different implementation of the Fast Fourier Transform
- Time Delay Estimation kernel
- Matrix Multiplication kernel

	Number of Buffers	Number of Actores	Number of Time Steps	Base- line	Coarse- Grain	Fine- Grain (Best Case)	Fine- Grain (Worst Case)	Compile Time with GA in Sec.	Optimal Solution by ILP
Bitonic Sort	119	214	340	1152	96	48	64	91	32
Insertion Sort	8	9	263	1024	256	128	128	6	128
FFT2	22	24	446	3072	640	384	384	10	~
FFT3	38	64	175	960	192	72	96	11	64
TDE	48	51	17204	77120	23168	11776	23040	510	~
Matrix Mult.	10	21	2712	5000	4000	2000	2000	13	~



Improvement of coarse-grain and fine-grain methods compared to the baseline.



Improvement in all fine-grain cases: GA worst case, GA best case, and ILP, compared to the coarse-grain method

#### Conclusions

- Visualization of buffers transforms the allocation problem into packing of complex polygons
- Fine-grain analysis vs. conventional coarse-grain live range analysis: dramatic improvements
- The benefits of this approach outweighs the reasonable increase in static analysis latency for a large class of resource-constrained embedded systems.